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Efficient Transmission in Multi-Antenna Two-Way AF Relaying Networks

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Abstract—In this paper, an efficient transmission scheme, termed the joint antenna selection and data exchange (AS-DE) scheme, is proposed for a two-way amplify-and-forward relaying network, where two single-antenna source terminals exchange information via a multi-antenna relay station. For the proposed scheme, the best antenna at the relay for each source terminal is first selected separately, following the max-max scheme. Then, from the set of the previously selected antennas, either one antenna is selected, in a similar fashion as well-known max-min and max-sum schemes, or two antennas exchange their respective received signals, which are then coded, amplified and broadcasted to the source and destination terminals. Tight lower and upper bounds on the outage probability (OP) for the proposed scheme have been derived assuming independent and identically distributed Rayleigh fading channels. Furthermore, our analysis reveals that the proposed joint AS-DE scheme can achieve full diversity. Finally, it is shown that under the same resource constraints, i.e., in terms of the number of the utilized time slots and transmit power, the proposed joint AS-DE scheme outperforms the max-min, the max-sum and the max-max schemes. Extensive numerical results accompanied with computer simulations, are further provided to validate the developed analytical results.

Index Terms—Two-way relaying networks, outage probabilities, antenna selection, max-min, max-sum, max-max.

I. INTRODUCTION

Recently, two-way relaying networks (TWRNs) have been envisioned as a promising transmission technology to significantly improve the reliability and transmission rate of wireless systems [1], [2]. The performance of TWRNs can be further improved by integrating multiple-input multiple-output (MIMO) transmission technology [3]–[5]. Antenna selection (AS), i.e., optimally choosing a subset of the available antennas, is an attractive low-cost and low-complexity technique, but still retains many of the advantages of conventional MIMO systems [6]. In the open technical literature, three antenna selection schemes for MIMO amplify-and-forward (AF) and decode-and-forward (DF) TWRNs have been proposed, namely the max-min [7], [8], the max-sum [9] and the max-max schemes [10], [11].

The performance achieved by such schemes has been assessed in several past research works. For example, the outage probability (OP) performance of the max-min and the max-sum schemes has been evaluated in [7]–[9]. These works have shown that both schemes can achieve full diversity. In [10], antenna selection in a DF relaying network based on the max-max scheme was investigated, assuming that decoding at the relay is error-free. In [11], the so-called double-max scheme was proposed. In that work, relay selection based on the max-max scheme was addressed, assuming the use of an error-free decoding relay.

Motivation: For the purpose of illustration, consider two single-antenna sources $T_1$ and $T_2$ exchanging information via a relay station $R$ which is equipped with $N = 3$ antennas, denoted by antenna $R_1$, antenna $R_2$ and antenna $R_3$, respectively. For example, let the channel gains from $T_1$ and $T_2$ to $R$ at a given time instant be $h = \{h_1, h_2, h_3\} = \{0.35, 0.46, 0.59\}$ and $g = \{g_1, g_2, g_3\} = \{0.72, 0.54, 0.32\}$, respectively. According to the max-min scheme, the best antenna at the relay is selected to maximize the end-to-end signal-to-noise ratio of the worse source [7], [8]. In this example, the antenna antenna $R_2$ will be chosen with $h_2 = 0.46$ and $g_2 = 0.54$. However, it can be observed that the links having the largest channel gains, i.e., $h_3 = 0.59$ and $g_1 = 0.72$, have not been utilized.

When the max-sum scheme is utilized, the best antenna at the relay is selected to maximize the sum-rate [9]. In the considered test case, the antenna $R_1$ will be chosen with $h_1 = 0.35$ and $g_1 = 0.72$. However, as it can be observed this method does not exploit the channel coefficient $h_3 = 0.59$, i.e., the maximum channel gain in all $h_i s$, $i \in \{1, 2, 3\}$.
The max-max scheme selects at a given time instance either one or two antennas at the relay, corresponding to the maximum channel coefficients [10], [11]. If the antenna indices are the same, one antenna is selected, otherwise two antennas are selected. In the previously described example, two antennas are selected, namely the antenna $R_1$ and antenna $R_3$ corresponding to the links with $h_3 = 0.59$ and $g_1 = 0.72$. Consider, however, the following data transmission scenario from $T_1$ to $T_2$. Specifically, assume that information flows from the links $T_1 \rightarrow R_1$, $R_1 \rightarrow T_2$ and $T_1 \rightarrow R_3$, $R_3 \rightarrow T_2$, characterized by channel gains $h_1 = 0.35$, $g_1 = 0.72$ and $h_3 = 0.59$, $g_3 = 0.32$, respectively. As can be observed, during the transmission through the antenna $R_1$, link $T_1 \rightarrow R_1$ experiences the worse channel conditions since $h_1$ is the minimum channel coefficient in $h_i$s. On the other hand, link $R_1 \rightarrow T_2$ experiences the best channel conditions because $g_1$ is the maximum channel coefficient in $g_i$s. Similar findings can be found when transmission through the links $T_1 \rightarrow R_3$, $R_3 \rightarrow T_2$ is considered. In such scenarios, the combinations of “small-maximum” and “maximum-small” channel coefficients result in a small received end-to-end (e2e) SNR at $T_2$. Note that when one antenna is selected, i.e., when the selected antennas’ indices are identical, data transmission will exploit the best links in an optimal way. In such a case, the max-max scheme exhibits the best performance. However, this is a small probability event.

Motivated by this key observation, in this paper, an efficient transmission scheme which can exploit the unutilized links characterized by the best channel coefficients, termed the joint antenna selection and data exchange (AS-DE) scheme, is proposed for multi-antenna AF TWRNs. The key idea in the joint AS-DE scheme is to combine max-max antenna selection scheme along with data exchange to transmit data through the links characterized by “maximum-maximum” and “small-small” channel coefficients. Consequently, the joint AS-DE scheme outperforms the max-min, the max-sum and the max-max schemes because its e2e SNR is significantly larger than that achieved by the aforementioned AS schemes. It should be emphasized that the previously reported works on the max-max scheme, such as those presented in [10], [11], ignore the possible transmission error due to the aforementioned “small-maximum” and “maximum-small” combinations of channel coefficients, since they consider DF relaying networks and assume decoding at the relay is error-free.

The performance of the joint AS-DE scheme is assessed by deriving the tight upper and lower bounds on the e2e OP, assuming Rayleigh fading conditions. The tightness of the newly derived bounds is verified by means of computer simulation. Extensive numerical results are further presented revealing that the joint AS-DE scheme can achieve full diversity. In addition, it is shown that under the same resource consumption constraints, such as in terms of the utilized time slots and transmit power, the joint AS-DE scheme also outperforms the existing max-min, max-sum and max-max schemes.

The remainder of this paper is organized as follows: Section II presents the system model and the joint AS-DE scheme. Section III investigates the OP and diversity gain performance for the joint AS-DE scheme. Numerical and simulation results are presented in Section IV. Finally, Section V concludes the paper.

Notation: $E\{\cdot\}$ and $I_M$ denote the expectation operation and an $M \times M$ identity matrix, respectively. $K_v(\cdot)$ and $Ei(\cdot)$ denote the $v$ order modified Bessel function of the second kind [12, Eq. (8.407)] and the Exponential integral function [12, Eq. (8.211)], respectively. The notations $CN(0,\sigma^2)$, $f_X(\cdot)$ and $F_X(\cdot)$ represent a circularly symmetric complex Gaussian random variable (RV) with zero mean and variance $\sigma^2$, the probability density function (PDF) and cumulative distribution function (CDF) of RV $X$, respectively. $\Pr(\cdot)$ returns the probability.

II. SYSTEM MODEL AND THE PROPOSED JOINT AS-DE SCHEME

In this section, the system model and the joint AS-DE scheme are introduced.

A. System Model

Consider a TWRN, where two single-antenna source terminals $T_1$ and $T_2$ exchange information by using an AF relay station $R$ equipped with $N \geq 2$ antennas. Assume that all the links experience independent and identically distributed (i.i.d.) Rayleigh fading, following $CN(0,\sigma)$, and channels are reciprocal. Assume that the $i$-th antenna $R_i$ is selected to help the communication between $T_1$ and $T_2$. The whole communication takes place in two times slots. In the first time slot, $T_1$ and $T_2$ transmit their signals to $R$. The received signal at the antenna $R_i$ after $M$ successive symbol durations can be written as

$$y_i = \sqrt{P}h_is_1 + \sqrt{P}g_is_2 + n_i,$$

where $s_j = [s_{j1}, \ldots, s_{jM}]^T$, $j = 1, 2$, denotes the transmitted symbol of $T_j$ with $E[s_j s_j^H] = I_M$, $P$ is the transmit power of $T_j$, $h_i$ and $g_i$ denote the channel coefficients between $T_1$ and antenna $R_i$, and between $T_2$ and the antenna $R_i$, respectively, $n_i \sim CN(0,N_0I_M)$ represents additive gaussian white noise (AWGN) at the antenna $R_i$.

In the second time slot, the selected antenna $R_i$ amplifies its received signal with gain $\alpha$ and then broadcasts it to $T_j$. The received signals at $T_1$ and $T_2$ are given by

$$y_{T_1} = \sqrt{P_r}h_i\alpha y_i + n_{T_1},$$

and

$$y_{T_2} = \sqrt{P_r}g_i\alpha y_i + n_{T_2},$$

respectively, where $P_r$ denotes the transmit power of the $i$-th antenna at $R$ and $n_{T_j} \sim CN(0,N_0I_M)$ is AWGN at terminal $T_j$. Assuming that fixed gain relaying is used, the amplification factor is expressed as $[1]$, $[13]$,

$$\alpha = \sqrt{\frac{1}{2P\Omega + N_0}}.$$  

After the self-interference cancellation is performed, assuming $P_r = 2P$, the received SNR at $T_1$ and $T_2$ via the help of the antenna $R_i$ is given as $[13]$,

$$\gamma_{T_1} = \frac{2\gamma_{T_1}^r + 1}{2\gamma_{T_1}^r + c},$$

and

$$\gamma_{T_2} = \frac{2\gamma_{T_2}^r + 1}{2\gamma_{T_2}^r + c},$$

(3)
where $\gamma_i^l = P|h_i|^2/N_0$, $\gamma_i^r = P|g_i|^2/N_0$, $\gamma = PQ/N_0$, and $c = 2\gamma + 1$.

In the following, the three conventional AS schemes, i.e., the max-min, the max-sum and the max-max schemes, are introduced.

- In the max-min scheme, the $i^*-\text{th}$ antenna is selected according to [7], [8],
  \[ i^* = \arg \max_{1 \leq i \leq N} \min(\gamma_i^T, \gamma_i^G). \]

- In the max-sum scheme, the selected antenna $i^*$ follows [9]
  \[ i^* = \arg \max_{1 \leq i \leq N} (1 + \gamma_i^T)(1 + \gamma_i^G). \]

From (4) and (5), it can be observed that only one antenna can be selected for relaying between $T_1$ and $T_2$ in both the max-min and max-sum schemes.

- In the max-max scheme, the $l^*$-th and $r^*$-th antennas are selected according to [10], [11],
  \[ l^* = \arg \max_{1 \leq i \leq N} h_i, \quad \text{and} \quad r^* = \arg \max_{1 \leq j \leq N} g_j. \]

From (6), it can be observed that if the antenna indices $l^* = r^*$, only one antenna can be used for relaying between $T_1$ and $T_2$ in the max-max scheme. Otherwise, two antennas can be used. Besides, it can be seen that the selected $l^*$-th and $r^*$-th antennas have the largest channel gain to $T_1$ and $T_2$, respectively.

**B. The Joint AS-DE Scheme**

The proposed joint AS-DE scheme includes two procedures, i.e., antenna selection and data exchange. Antenna selection is performed based on the max-max scheme in a similar fashion as in (6). During the data exchange phase, the $l^*$-th antenna at $R$ transmits its signal $y_{l^*}$ to the $r^*$-th antenna at $R$, and the $r^*$-th antenna at $R$ transmits its signal $y_{r^*}$ to the $l^*$-th antenna at $R$. We note that data exchange is started only when two antennas are selected, i.e., the antenna indices $l^* \neq r^*$. The whole communication takes place in two time slots.

Let us now rearrange $h_i$ and $g_i$, $i = 1, \ldots, N$, in an ascending order. We define the channel coefficients $h_{(i)}$ and $g_{(i)}$, respectively, such that $h_{(1)} \leq h_{(2)} \leq \cdots \leq h_{(N)} = h$, and $g_{(1)} \leq g_{(2)} \leq \cdots \leq g_{(N)} = g$. Note that $h$ and $g$ are the $N$-th largest channel gain in $h_i$ and $g_i$, respectively. In the following, data transmission in the joint AS-DE scheme is presented when either one or two antennas are selected.

1) **If One Antenna is Selected:** In this case, the antenna indices are the same, i.e., $l^* = r^*$. In the first time slot, $T_1$ and $T_2$ broadcast their information to $R$. In the second time slot, the selected antenna amplifies its received signal in the first time slot by a gain $\alpha$, and broadcasts to $T_1$ and $T_2$ with full power $P_r = 2P_r$.

Similar to (3), the received SNR at $T_2$ can be obtained as
\[ \gamma_{1,2} = \frac{2\gamma_1^l\gamma_1^r}{2\gamma_1^r + c}. \]

2) **If Two Antennas are Selected:** In this case, the antenna indices are different, i.e., $l^* \neq r^*$. In the first time slot, the received signals at $R$ via the antennas $l^*$ and $r^*$, which have been selected based on (6), are given by
\[ y_{l^*} = \sqrt{P}h_{l^*}s_1 + \sqrt{P}g_{l^*}s_2 + n_{l^*}, \]
and
\[ y_{r^*} = \sqrt{P}h_{r^*}s_1 + \sqrt{P}g_{r^*}s_2 + n_{r^*}, \]
respectively. The $n_{l^*}$, $n_{r^*} \sim CN(0, N_0I_M)$ represent the AWGN at the $l^*$-th and $r^*$-th antennas, respectively.

Following [13], the durations of both time slots are considered to be the same. Data exchange between the $l^*$-th and $r^*$-th antennas occurs in the second time slot. Then, the selected $l^*$-th antenna transmits its signal $y_{l^*}$ to the $r^*$-th antenna, and the $r^*$-th antenna transmits its signal $y_{r^*}$ to the $l^*$-th antenna. In the same time slot, the $l^*$-th and $r^*$-th antennas process $y_{l^*}$ and $y_{r^*}$ to generate the space time coded symbol $x_{l^*}$ and $x_{r^*}$, respectively. The transmitted signals $x_{l^*}$ and $x_{r^*}$ are designed to be linear functions of $y_{r^*}$ and $y_{l^*}$ and their conjugates, namely [13]-[16],
\[ x_{l^*} = A_{l^*}y_{r^*} + B_{l^*}y_{r^*}^*, \quad \text{and} \quad x_{r^*} = A_{r^*}y_{l^*} + B_{r^*}y_{l^*}^*, \]
where $A_p$ and $B_p$, $p \in \{l^*, r^*\}$, are $M \times M$ precoding matrices, designed using guidelines for the construction of distributed space-time coding schemes, and $y_p^*$ denotes the conjugate of $y_p$. For simplicity, in this paper, we consider $M = 2$, and the orthogonal matrices are used at the two selected antennas as in [16], namely
\[ A_{l^*} = I_2, B_{l^*} = 0_2, A_{r^*} = 0_2, B_{r^*} = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}. \]

Therefore, (8) becomes
\[ x_{l^*} = A_{l^*} \left( \sqrt{P}h_{r^*}s_1 + \sqrt{P}g_{r^*}s_2 + n_{r^*} \right), \]
and
\[ x_{r^*} = B_{r^*} \left( \sqrt{P}h_{l^*}s_1 + \sqrt{P}g_{l^*}s_2 + n_{l^*} \right)^*. \]

Then, the $l^*$-th and $r^*$-th antennas broadcast $x_{l^*}$ and $x_{r^*}$ after amplification, respectively, each with half power $P_r/2 = P_r$. Let $\hat{y}_{T_2}$ denote the received signals after self-interference cancellation at $T_2$, given by
\[ \hat{y}_{T_2} = \sqrt{\frac{P_r}{2}} g_{l^*}\alpha' B_{r^*} \left( \sqrt{P}h_{l^*}s_1 + n_{l^*} \right)^* \]
\[ + \sqrt{\frac{P_r}{2}} g_{r^*}\alpha' A_{l^*} \left( \sqrt{P}h_{r^*}s_1 + n_{r^*} \right) + n_{T_2}. \]

The average transmit power at each antenna is constrained to
\[ \mathbb{E} \left\{ ||\alpha' x_p||_F^2 \right\} = 1. \]
When orthogonal matrices in (9) are employed, $\alpha' = \alpha \sqrt{2}$, and the received SNR at $T_2$ can be obtained, based on (10), as follows

$$\gamma_{ASDE}^{T_2} = \frac{\alpha'^2 \mathbb{P}|h_{1r}|^2|g_{1r}|^2||B_{rT}||^2 + \alpha'^2 \mathbb{P}|h_{2r}|^2|g_{2r}|^2||A_{rT}||^2}{\alpha'^2 \mathbb{P}|g_{1r}|^2||B_{rT}||^2N_0 + \alpha'^2 \mathbb{P}|g_{2r}|^2||A_{rT}||^2N_0 + 2N_0}.$$  \hspace{1cm} (12)

Furthermore, (12) can be re-expressed as

$$\gamma_{ASDE}^{T_2} = \frac{\gamma_{(N)}\gamma_{(I)} + \gamma_{(q)}\gamma_{(w)}}{\gamma_{(q)} + \gamma_{(N)} + c}.$$  \hspace{1cm} (13)

Eq. (13) can be upper- and lower-bounded as

$$\gamma_{ASDE}^{T_2,ub} \geq \gamma_{ASDE}^{T_2} \geq \gamma_{ASDE}^{T_2,lb},$$  \hspace{1cm} (14)

where

$$\gamma_{ASDE}^{T_2,lb} = \frac{\gamma_{(N)}\gamma_{(I)} + \gamma_{(q)}\gamma_{(w)}}{\gamma_{(q)} + \gamma_{(N)} + c},$$  \hspace{1cm} (15)

$$\gamma_{ASDE}^{T_2,ub} = \frac{\gamma_{(N)}\gamma_{(I)} + \gamma_{(q)}\gamma_{(w)}}{\gamma_{(q)} + \gamma_{(N)} + c}.$$  \hspace{1cm} (16)

To compare the joint AS-DE scheme with the max-max scheme, here, we present the received SNR at $T_2$ in the max-max scheme. In the max-max scheme, since it does not utilize the “data exchange”, in the second time slot, the transmitted signal $x_r'$ at the $r$-th antenna and the transmitted signal $x_r^*$ at the $r^*$-th antenna are given as

$$x_r' = A_r \left( \sqrt{P} h_{rT} s_3 + \sqrt{P} g_{rT} s_2 + n_r \right),$$

and

$$x_r^* = B_r \left( \sqrt{P} h_{rT} s_3 + \sqrt{P} g_{rT} s_2 + n_r \right)^*,$$

respectively.

Following similar arguments as to (13), the received SNR at $T_2$ in the max-max scheme can be obtained as

$$\gamma_{MaxMax}^{T_2} = \frac{\gamma_{(q)}\gamma_{(I)} + \gamma_{(N)}\gamma_{(w)}}{\gamma_{(q)} + \gamma_{(N)} + c}.$$  \hspace{1cm} (17)

**Remark 1:** Comparing (13) with (17), it can be seen that the difference between the two SNR results is in their numerators. In the numerator of (13), there exists the “maximum-maximum”, i.e., $\gamma_{(N)}\gamma_{(I)}$, and “small-small”, i.e., $\gamma_{(q)}\gamma_{(w)}$, channel coefficient combinations for the joint AS-DE scheme, but “small-maximum” and “maximum-small”, i.e., $\gamma_{(q)}\gamma_{(N)}$ and $\gamma_{(N)}\gamma_{(w)}$, link combinations in (17) for the max-max scheme. This is because both $\gamma_{(N)}$ and $\gamma_{(I)}$ are the maximal effective channel gain, due to the fact that $g_r$ and $h_r$ are the $N$-th maximum among $g$s and $h$s; and, $\gamma_{(q)}$ and $\gamma_{(w)}$ are small, due to the fact that $g_r$ and $h_r$ are the $q$-th maximum and $w$-th maximum in $g$s and $h$s, respectively. The “maximum-maximum” and “small-small” link combinations in the joint AS-DE scheme, result in a larger received SNR than in the case of the max-max scheme, where the “small-maximum” and “maximum-small” link combinations are used. Although the max-max scheme exploits the best links with the largest channel gains, it does not utilize them in the best manner. Recalling the example described in the introduction section, one can find that the max-min and the max-sum schemes may not exploit the best links. However, in the joint AS-DE scheme data transmission from $T_1$ to $T_2$ uses the links $T_1 \to R_3$, $R_1 \to T_2$ and $T_1 \to R_1$, $R_3 \to T_2$, having gains $h_3 = 0.59$, $g_1 = 0.72$ and $h_1 = 0.35$, $g_3 = 0.32$, respectively. From this and (13), we conclude that the joint AS-DE scheme exploits the strong channel links in the best manner, by utilizing “maximum-maximum” and “small-small” channel coefficients combinations. Because of this, it outperforms the max-max, the max-min and the max-sum schemes.

**Remark 2:** Despite the fact that the joint AS-DE scheme outperforms the conventional AS schemes, it requires a second RF chain for its practical implementation. Moreover, its baseband implementation is more complicated than the one of conventional schemes, as it requires data exchange between the selected antennas. Recently, novel MIMO transmission schemes have been reported, such as spatial modulation and MIMO electronically steerable passive array radiator (ESPAR) [17]–[19]. Such schemes can minimize complexity and the costs while attaining the advantages of the MIMO system.

**Remark 3:** Hereafter some issues regarding the implementation of the data exchange phase of the joint AS-DE scheme are discussed. Such a scheme can be implemented in an efficient manner in baseband, by employing digital hardware, instead of exchanging analog signals between antennas. Specialized devices, such as digital signal processors (DSP) or field programmable gate arrays (FPGA) can be used to this purpose. Such devices are equipped with specialized direct memory access (DMA) controllers, thus rendering them capable of transferring large amounts of data, stored in buffers. Sophisticated techniques, such as multiple buffering, can be also employed to increase the efficiency of data transmission. Data exchange between two antennas can be implemented without significant computational complexity by exchanging the contents of their corresponding buffers. For a given hardware platform, one can perform such a task in an optimized way, i.e., by minimizing the number of the required clock cycles.

### III. Performance Analysis

In this section, the OP performance at $T_2$ will be analyzed. The performance analysis at $T_1$ can be obtained in a similar fashion and thus mathematical derivations are omitted for brevity. For writing simplicity, some definitions are given as follows.

**Definition 1:**

$$\sum_{1,k_1} \sum_{k_i=0}^{N} \binom{N}{k_i} (-1)^{k_i} \sum_{2,k_i} \sum_{k_i=0}^{N} \binom{N}{k_i} (-1)^{k_i+1}k_i,$$

and

$$\sum_{3,k_1,k_2} \frac{1}{\pi^2} \sum_{k_{i}=0}^{N-1} \sum_{k_{j}=0}^{N-1-q} \binom{N}{k_{i}} \binom{q-1}{k_{j}} \left( \frac{1}{N - q} \right)^{k_{i}+k_{j}}.$$
The distributions of $\gamma_{1,2}$ and $\gamma_{ASDE}$ will be firstly presented in the following theorems, which lay the foundation for performance analysis.

A. Distribution of the received SNR

**Theorem 1:** When one antenna is selected for relaying between $T_1$ and $T_2$, i.e., the antenna indices $l^* = r^*$, the CDF of the received SNR at terminal $T_2$, i.e., $\gamma_{1,2}$, can be expressed in closed-form as

$$F_{\gamma_{1,2}}(z) = 1 + \sum_{1,k_1 \neq 0} \sum_{2,k_2 \neq 0} \frac{2k_1 k_2}{k_2} e^{-\frac{k_1}{k_2}} K_1 \left( \sqrt{\frac{2k_1 k_2 z}{\gamma}} \right).$$

**Proof:** See Appendix A.

**Theorem 2:** When two antennas are selected for relaying between $T_1$ and $T_2$, i.e., the antenna indices $l^* \neq r^*$, the CDF of the upper bound on $\gamma_{2,2}$, i.e., $F_{\gamma_{ASDE}}(z)$, is given as

$$F_{\gamma_{ASDE}}(z) = \begin{cases} \mathcal{L}_1, & a = b \\ \mathcal{L}_2, & a \neq b \end{cases}$$


where

$$a = \frac{k_2 + 1}{\gamma}, \quad b = \frac{N + k_1 - q - k_2}{\gamma},$$

$$\mathcal{L}_1 = \sum_{3,k_1,k_2} \left[ \frac{1}{a(a + b)} + \sum_{1,k_3 \neq 0} \frac{k_3 c z}{a \gamma} e^{-\frac{k_3}{a \gamma}} K_2 \left( 2 \sqrt{\frac{k_3 a c z}{\gamma}} \right) \right]$$

and

$$\mathcal{L}_2 = \sum_{3,k_1,k_2} \left[ \frac{1}{a(a + b)} + \sum_{1,k_3 \neq 0} \frac{2 c e^{-\frac{k_3}{a \gamma}}}{b - a} K_1 \left( 2 \sqrt{\frac{k_3 a c z}{\gamma}} \right) - \sqrt{\frac{2k_3 z}{(a + b) c \gamma}} K_1 \left( \sqrt{\frac{2k_3(a + b) c z}{\gamma}} \right) \right].$$

**Proof:** See Appendix B.

**Theorem 3:** When two antennas are selected, i.e., $l^* \neq r^*$, the CDF of the lower bound on $\gamma_{2,2}$, i.e., $F_{\gamma_{ASDE}}(z)$, is given as

$$F_{\gamma_{ASDE}}(z) = \sum_{3,k_1,k_2,k_3} \left( \int_{0}^{1} e^{-\frac{k_3}{a \gamma}} A_1 d\zeta + \int_{0}^{0.5} e^{-\frac{2k_3}{a \gamma}} A_2 d\zeta \right)$$

where

$$A_1 = -\frac{1 + \frac{a c}{\gamma} + \frac{b c}{\gamma}}{(\zeta - 1)(\zeta) + b^2} e^{-\frac{a c}{\gamma}},$$

and

$$A_2 = -\frac{e^{-\frac{a c}{\gamma}}}{(\zeta - 1)(\zeta) + b^2} \left( -1 + \frac{a c}{\gamma} + \frac{b c}{\gamma} \right).$$

**Proof:** See Appendix C.

**Remark 4:** Theorem 3 involves the computation of the integrals with integrands composed of elementary functions. Although such integrals are not in closed-form, they can be easily evaluated numerically by employing standard techniques available in the most common mathematical software packages, such as Matlab, Maple, or Mathematica.

B. Outage Probability

The OP is defined as the probability that the instantaneous SNR falls below a given threshold $\gamma_{th}$, i.e.,

$$P_{out}(\gamma_{th}) = \Pr[\gamma < \gamma_{th}] = F_{\gamma}(\gamma_{th}).$$

We note that the OP at $T_j$, $j = 1, 2$ is given by

$$P_{out}(T_j) = \Pr[\gamma < \gamma_{th}]$$

The OP results are presented in the following corollaries.

**Corollary 1:** In the joint AS-DE scheme, the tight upper and lower bounds on the OP at $T_2$ can be calculated as

$$P_{out,ub}(\gamma_{th}) = p_N F_{\gamma_{ASDE}}(\gamma_{th}) + p_N \sum_{q=1}^{N-1} F_{\gamma_{ASDE}}(\gamma_{th}),$$

and

$$P_{out,lb}(\gamma_{th}) = p_N F_{\gamma_{ASDE}}(\gamma_{th}) + p_N \sum_{q=1}^{N-1} F_{\gamma_{ASDE}}(\gamma_{th}),$$

respectively, where $p_N = p_N’ = 1/N$, $F_{\gamma_{ASDE}}(\gamma_{th})$ and $F_{\gamma_{ASDE}}(\gamma_{th})$ are presented in Theorem 1, Theorem 2 and Theorem 3, respectively.

**Remark 5:** The probabilities $p_N$ in (24) and (25) are for the event that one antenna is selected. From (6), it can be seen that among the $N \times N$ pairs $(h_i, g_j)$, there are $N$ pairs $(h_i, g_j)$ with $l^* = r^* = 1, \cdots, N$. We note that each pair is selected with the same probability, since we assume all $h_i$s and $g_j$s are i.i.d. distributed. Therefore, the probability that one antenna is selected, is $p_N = N/(N \times N) = 1/N$. In (24) and (25), $p_N’$ is equal to $(1 - 1/N)/(N - 1)$, where $1 - 1/N$ is the probability that two different antennas have been selected, i.e., $l^* \neq r^*$, and $1/(N - 1)$ is the probability that $q$ takes a specific value in $\{1, \cdots, N - 1\}$.

C. Diversity Order

**Corollary 2:** The proposed joint AS-DE scheme can achieve full diversity, i.e., the diversity order is $N$.

**Proof:** See Appendix D.

IV. SIMULATION AND NUMERICAL RESULTS

In this section, computer simulations are carried out to demonstrate the performance of the joint AS-DE scheme with $\gamma_{th} = 3$, $P = 1$ and $P_r = 2P = 2$ in all figures, $\Omega = 1$ in Figs. 1-3, and $\Omega = 1/(1 - d)^{-3}$ in Fig. 4 where $d$ and $3$ denote the distance between $T_1$ and $R$ and the path-loss exponent, respectively. In addition, numerical results obtained from Corollary 1 are also used to show the accuracy of the developed analytical results. We note that in order to guarantee comparison fairness, the same power consumption used by the joint AS-DE scheme as in the max-min, the max-sum and the max-max schemes is considered. For the max-min and max-sum schemes which select a single antenna at $R$,
the whole transmit power at $R$ is $P_r = 2$. For the proposed AS-DE and max-max schemes, if one antenna is selected, the transmit power of the selected antenna at $R$ is $P_r = 2$; otherwise, each selected antenna broadcasts the signal with power $P_r/2 = 1$, indicating that the whole transmit power at $R$ is $P_r + P_r/2 = 2$.

Fig. 1 illustrates the Monte-Carlo simulation results on the OP performance of the joint AS-DE scheme in comparison to the max-min, max-sum, max-max schemes versus $\tau$ under $N = 2, 4$. It clearly illustrates that under arbitrary $N$ and $\tau$, our proposed joint AS-DE scheme performs much better than the other three AS schemes. For example, when $N = 4$ and at $10^{-4}$ OP, the joint AS-DE scheme provides a nearly 3 dB gain over those of the max-min and max-sum schemes, about a 10 dB gain over that of the max-max scheme. This result is expected because the joint AS-DE scheme utilizes the "maximum-maximum" and "small-small" channel coefficients combinations, resulting in a larger received SNR. Furthermore, it can also be seen that the joint AS-DE scheme can achieve full diversity as the max-min, max-sum schemes.

Fig. 2 compares the OP performance of the joint AS-DE scheme with the max-min, max-sum and max-max schemes versus the relay antenna number $N$ with $\tau = 10$ dB. It can be also clearly seen that the joint AS-DE scheme outperforms the other three schemes under an arbitrary $N$.Besides, as the number of relay antennas $N$ increases, the SNR gain that the joint AS-DE scheme achieves over the other three schemes counterpart is further enlarged.

In Fig. 3, the developed analytical results presented in Corollary 1 for the joint AS-DE schemes are compared to the simulation results with $N = 2, 3$ and 4. As can be
observed from the figure, under an arbitrary \( N \), the upper and lower bounds on OP are quite close to the simulation counterparts which verifies our analysis. Furthermore, the lower and upper bound outage curves verify our diversity order analysis, indicating that the joint AS-DE scheme can achieve full diversity.

In Fig. 4, the impact of the relay location on OP for \( T_2 \) is studied in the joint AS-DE scheme. Specifically, the OP is plotted against the distance \( d \) between \( T_1 \) and \( R \) by modeling the path-loss dependent parameters \( \Omega = 1/(1 - d)^{-3} \). As can be observed, the OP performance at \( T_2 \) improves as the relay station gets close to \( T_2 \). Besides, the lower and upper bounds we derived are quite tight under an arbitrary \( d \) indicating the accuracy of our analysis.

V. CONCLUSION

In this paper, we proposed an efficient transmission scheme for multi-antenna AF TWRNs, termed as the joint AS-DE scheme. Particularly, the joint AS-DE scheme utilized the antenna selection criterion in the max-max scheme along with data exchange to transmit data through the links characterized by “maximum-maximum” and “small-small” channel coefficient combinations. We presented the tight lower and upper bounds on the OP for the proposed scheme. Furthermore, our analysis revealed that the joint AS-DE scheme can achieve full diversity. Finally, analysis and simulation results showed that under the same time slots and power consumption, the joint AS-DE scheme outperforms the existing schemes, i.e., the max-min, max-sum and max-max ones. For example, when \( N = 4 \) and at \( 10^{-4} \) OP, the joint AS-DE scheme provides a nearly 3 dB gain when compared with the max-min and max-sum schemes, and about a 10 dB gain when compared with the max-max scheme.

APPENDIX A

PROOF OF THEOREM 1

Since all links experience i.i.d. Rayleigh fading, the PDF and CDF of the instantaneous SNR of any links, \( \gamma_1 \) or \( \gamma_1^r \) follow that

\[
 f_{\gamma_1}(x) = \frac{1}{x} e^{-\frac{x}{\theta}}, \quad F_{\gamma_1}(x) = 1 - e^{-\frac{x}{\theta}}.
\]

Based on the order statistics in [20], we have

\[
 F_{\gamma_1(N)}(x) = (F_{\gamma_1}(x))^N = \sum_{1,k_1} e^{-\frac{k_1 x}{\theta}}. \quad \text{(A-1)}
\]

The PDF of \( \gamma_1(N) \) can be obtained as,

\[
 f_{\gamma_1(N)}(x) = \sum_{2,k_2} e^{-\frac{2k_2 x}{\theta}}.
\]

Similarly, we have

\[
 f_{\gamma_1(N)}(y) = \sum_{2,k_2} e^{-\frac{2k_2 y}{\theta}}.
\]

From (7), \( F_{\gamma_1,T_2}(z) \) can be expressed as

\[
 F_{\gamma_1,T_2}(z) = \int_0^\infty \Pr \left( \gamma_1(N) \leq z + \frac{2cz}{2y} \right) f_{\gamma_1(N)}(y)dy.
\]

Utilizing [12, Eq.(3.471.9)], Theorem 1 can be achieved.

APPENDIX B

PROOF OF THEOREM 2

We will firstly study the distribution of \( \theta = \gamma_q + \gamma_q^r \), and then the distribution of \( u = \theta / (\theta + c) \). Finally, the distribution of \( \gamma_{ASDE,T_2,ub} = \gamma_l(N)u \) will be obtained.

Now, let’s study the CDF of \( \theta = \gamma_q + \gamma_q^r \). Based on the order statistics in [20], the joint PDF of \( \gamma_q \) and \( \gamma_q^r \), \( 1 \leq q < N \), is

\[
 f_{\gamma_q,\gamma_q^r}(s,v) = \sum_{3,k_1,k_2} e^{-bs} e^{-av}, \quad \text{(B-1)}
\]

for \( 0 < s < v < \infty \).

Therefore, the CDF of \( \theta \), i.e., \( F_\theta(\theta) \), follows that

\[
 F_\theta(\theta) = \int_0^\theta \int_s^\infty f_{\gamma_q}(s,v)dvds = \sum_{3,k_1,k_2} \frac{1}{\theta^2} \int_0^{\theta/2} \left[ e^{-(a+b)s} - e^{-(a-b)s-a\theta} \right] ds. \quad \text{(B-2)}
\]

Utilizing [12, Eq. (2.311)], we have

\[
 F_\theta(\theta) = \begin{cases} 
 - \sum_{3,k_1,k_2} \frac{1}{\theta^2} \left[ a^2 e^{-a\theta} + \frac{e^{(a+b)\theta} - 1}{a+b} \right], & a-b = 0, \\
 \sum_{3,k_1,k_2} \frac{1}{\theta^2} \left[ \frac{e^{-(a-b)\theta}}{a-b} + \frac{e^{-(a+b)\theta}}{a+b} \right], & a-b \neq 0.
\end{cases} \quad \text{(B-3)}
\]

From \( u = \theta / (\theta + c) \), we have \( \theta = g(u) \), where \( g(u) = \frac{cu}{1+cu} \). Therefore, \( f_u(u) = f_\theta(g(u))g'(u) \), where \( g'(u) \) denotes the derivative of \( g(u) \).

Taking the derivative of (B-3), we can obtain the PDF of \( \theta \), i.e. \( f_\theta(\theta) \). And then, \( f_u(u) \) can be obtained as follows,

\[
 f_u(u) = \begin{cases} 
 - \sum_{3,k_1,k_2} \frac{e^{2cu}}{2(1-u)^2} e^{\frac{cu}{1-u}}, & a = b, \\
 \sum_{3,k_1,k_2} \frac{e^{cu}}{2(1-u)^2} \left( e^{\frac{cu}{1-u}} - e^{\frac{(a+b)cu}{2-a}} \right), & a \neq b, \\
 0, & u > 1.
\end{cases} \quad \text{(B-4)}
\]

The CDF of \( \gamma_{ASDE,T_2,ub} \) is given as follows

\[
 F_{\gamma_{ASDE,T_2,ub}}(z) = \int_{-\infty}^{\infty} F_{\gamma_l}(\frac{z}{u}) f_u(u)du. \quad \text{(B-5)}
\]

Substituting (A-1) and (B-4) into (B-5), with the aid of [12, Eq. (3.351.2)] and [12, Eq. (3.471.9)], Theorem 2 is deduced.

APPENDIX C

PROOF OF THEOREM 3

We firstly study the distribution of \( \varepsilon = \gamma_q(N) / (\gamma_q + \gamma_q^r + c) \), and then the distribution of \( \gamma_{ASDE,T_2,ub} = \varepsilon \gamma_q^l(N) \).

The CDF of \( \varepsilon \) is given as

\[
 F_\varepsilon(\zeta) = \Pr \left( \frac{\gamma_q(N)}{\gamma_q + \gamma_q^r + c} \leq \zeta \right). \quad \text{(C-1)}
\]
When $1/2 \leq \zeta \leq 1$, substituting (B-1) into (C-1), we have
\[
F_\zeta(\zeta) = \sum_{3,k_1,k_2} \int_0^\infty \int_0^{\frac{\zeta-1}{a+b}} e^{-a v} e^{-b s} \, dv \, ds
\]
\[
= \sum_{3,k_1,k_2} \frac{1}{a(a+b)} + \frac{(-1+\zeta) e^{-a+b}}{a+b(a-b)\zeta}.
\]

When $0 \leq \zeta \leq 1/2$, (C-1) becomes
\[
F_\zeta(\zeta) = \sum_{3,k_1,k_2} \int_{-\infty}^{\frac{\zeta-1}{a+b}} \int_{-\infty}^{\frac{\zeta-1}{a+b}} e^{-a v} e^{-b s} \, dv \, ds
\]
\[
= \sum_{3,k_1,k_2} \left( A_1 + \frac{(-1+\zeta) e^{-a+b}}{a(b-a)\zeta} - \frac{e^{-a+b}}{a(b-a)\zeta} \right).
\]

Therefore, the PDF of $\varepsilon$ can be obtained as,
\[
f_\varepsilon(\varepsilon) = \begin{cases} \sum \frac{A_1}{3,k_1,k_2}, & 1/2 < \zeta < 1, \\ \sum (A_1 + A_2), & 0 < \zeta \leq 1/2, \\ 0, & \text{others}. \end{cases}
\]

where $A_1$ and $A_2$ are given in (21) and (22), respectively. Therefore, the CDF of $\gamma_{\text{ASDE}}$, i.e., $F_{\gamma_{\text{ASDE}}}(\gamma_{th})$, is
\[
F_{\gamma_{\text{ASDE}}}(\gamma_{th}) = \int_0^\infty F_\zeta(\zeta) \frac{\gamma_{th}}{\zeta} f_\zeta(\zeta) d\zeta.
\]

Substituting (A-1) and (C-2) into (C-3), Theorem 3 can be reached.

APPENDIX D

PROOF OF COROLLARY 2

The following facts are utilized, i.e., $\lim_{x \rightarrow 0} e^{-x} = 1 - x$, $\lim_{x \rightarrow 0} K_1(x) = 1/x$ and $\lim_{x \rightarrow 0} K_2(x) = 2/x^2$.

Recalling Theorem 1, (18) can be re-expressed as,
\[
F_{\gamma_{1,\text{TH}}}(z) = 1 + \sum_{k_2=1}^N \frac{N}{k_2} (-1)^{k_2+1} \sum_{k_1=1}^N \frac{N}{k_1} (-1)^{k_1} e^{-b z}
\]
\[
= 1 + \left[ (1 - e^{-b z})^N - 1 \right].
\]

When $\gamma \rightarrow \infty$, we have
\[
F_{\gamma_{1,\text{TH}}}(z) \xrightarrow{\gamma \rightarrow \infty} \left( \frac{z}{\gamma} \right)^N.
\]

Utilizing (A-1), (C-3) can be rewritten as
\[
F_{\gamma_{1,\text{TH}}}(z) = \int_0^\infty (1 - e^{-b z})^N f_\zeta(\zeta) d\zeta
\]
\[
\xrightarrow{\gamma \rightarrow \infty} \left( \frac{z}{\gamma} \right)^N \int_0^\infty \frac{1}{\zeta^N} f_\zeta(\zeta) d\zeta.
\]

We note that $\int_0^\infty \frac{1}{\zeta^N} f_{\zeta}(\zeta) d\zeta$ is a constant, which is independent of $\gamma$.

Based on (D-1) and (D-2), at the high SNR regions, (24) can be asymptotically approximated by
\[
F_{\gamma_{1,\text{TH}}}(\gamma_{th}) \xrightarrow{\gamma \rightarrow \infty} \frac{1}{N} \sum_{q=1}^{N-1} \int_0^\infty \frac{1}{\zeta^N} f_{\zeta}(\zeta) d\zeta \left( \frac{\gamma_{th}}{\gamma} \right)^N.
\]

Recalling Theorem 2, $L_1$ in (19) can be deduced as, when $\gamma \rightarrow \infty$,
\[
L_1 \xrightarrow{\gamma \rightarrow \infty} \sum_{3,k_1,k_2} \frac{1}{a(a+b)} + \frac{1}{2a^2} \sum_{k_3=1}^N \left( N \right) (-1)^{k_3} e^{-b z}
\]
\[
= \sum_{3,k_1,k_2} \left\{ \frac{1}{a(a+b)} + \frac{1}{2a^2} \left[ (1-e^{-b z})^N - 1 \right] \right\}
\]
\[
= \sum_{3,k_1,k_2} \frac{1}{a(a+b)} \left( \frac{z}{\gamma} \right)^N.
\]

Similarly, $L_2$ in (19) can be obtained as follows, when $\gamma \rightarrow \infty$,
\[
L_2 \xrightarrow{\gamma \rightarrow \infty} \sum_{3,k_1,k_2} \frac{1}{a(a+b)} \left( \frac{z}{\gamma} \right)^N.
\]

Utilizing (D-4) and (D-5), when $\gamma \rightarrow \infty$, (19) is deduced as,
\[
F_{\gamma_{1,\text{TH}}}(z) \xrightarrow{\gamma \rightarrow \infty} \sum_{3,k_1,k_2} \frac{1}{a(a+b)} \left( \frac{z}{\gamma} \right)^N.
\]

From (D-1) and (D-6), at the high SNR regions, (25) can be asymptotically approximated by
\[
F_{\gamma_{1,\text{TH}}}(\gamma_{th}) \xrightarrow{\gamma \rightarrow \infty} \frac{1}{N} \sum_{q=1}^{N-1} \sum_{3,k_1,k_2} \frac{1}{a(a+b)} \left( \frac{\gamma_{th}}{\gamma} \right)^N.
\]

Finally, Corollary 2 is proved from (D-7) and (D-3).

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