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MO-MCS, a derivative-free algorithm for the multiobjective optimization of adsorption processes

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Abstract: Cyclic adsorption processes for gas separation, such as pressure and temperature swing adsorption (PSA and TSA), are non-stationary multi-column processes. Their design involves many degrees of freedom, which offers a very high flexibility while calling for a systematic and rigorous optimization approach. As an additional challenge, optimization of these separation processes involves multiple objectives, e.g. minimal energy demand and maximal productivity, which have to be pursued while fulfilling given process specifications, e.g. purity and recovery of the target components. This work extends the multilevel coordinate search (MCS), a well-known model-based derivative free algorithm, to constrained multiobjective problems. The algorithm, called MO-MCS, combines a built-in parallel computing strategy with the use of look-up tables with the goal of minimizing the computational time needed to determine the Pareto curve. The comparison with state-of-the-art optimizers indicates that MO-MCS shows better performance in terms of optimality, well spacing and extension of the Pareto curve. Afterwards, two industrially relevant case studies (TSA for CO₂ separation and PSA for H₂ and CO₂ co-production) are tackled to demonstrate the effectiveness of the algorithm as a tool to guide the design of adsorption processes.
1 Introduction

Adsorption-based separation technologies are often regarded as key players for advanced gas separation processes: solid sorbents are non-volatile, non-toxic, and offer the opportunity for a reduced energy penalty as compared to liquid sorbents, thanks to the lower sorption enthalpy. Cyclic adsorption processes have particularly attracted interest because of their capability in achieving high separation targets, flexibility to treat different compositions as feed, and simplicity of the equipment and sorbents. These processes are divided depending on the nature of the regeneration step: pressure swing adsorption (PSA) cycles exploit a pressure switch, while temperature swing adsorption (TSA) works with a temperature switch. Both PSA and TSA share the common principle, that the separation is achieved via using a solid medium undergoing alternating loading and regeneration phases. Accordingly, they are also hindered by very similar criticalities, especially: (i) handling of impurities, (ii) development of new, cost-effective and robust materials, (iii) design of suitable cycles, (iv) design of efficient bed geometries, including heat integration possibilities, and (v) enabling rigorous process optimization. The interplay among these challenges is especially tight when considering carbon capture and storage (CCS), a technology that is key for meeting the reduction of CO$_2$ emissions as required by the 1.5/2°C scenario. The present contribution aims at providing a new tool for the numerical optimization of cyclic adsorption processes, therefore tackling the last point of the list above.

The optimization of cyclic adsorption processes is inherently complex because of the dynamic and cyclic behaviour of such processes, and the tight relation between reactor design, cycle time steps and scheduling. Because of the dynamic operation of the adsorption bed, and the steep temperature and concentration profiles, the process simulation involves partial derivative equations (PDEs). Moreover, the complexity of the resolution is increased by the Cyclic Steady State condition (CSS) – i.e. cycle N+1 matches cycle N within the assigned tolerances.

Solution methods can be classified into two main categories: finite difference/elements methods, and shooting methods. In the first, the periodic CSS condition and the design constraints are solved together with the time-space discretized PDEs by a nonlinear equation solver (Nilchan and Pantelides). This leads to a very large set of algebraic equations. In the second (i.e. shooting methods), where the PDEs are converted into a set of differential algebraic equations (DAEs) and integrated over time until the adsorption cycle reaches CSS. The basic approach to handle the CSS condition is found by the successive substitution method, which consists in using the bed state at the end of the cycle as starting state for the subsequent cycle. While such approach mimics the behaviour of the real system, providing insights in the CSS convergence, it also features a linear convergence rate with considerable simulation times. More sophisticated approaches use sensitivity information from the solution of the DAEs to determine a new estimate of initial bed conditions using Newton-based algorithms which can achieve quadratic or superlinear convergence rates. On the other hand, these algorithms require extra computation time to calculate the Jacobian and may fail to converge for strongly nonlinear systems and complex bed models, because the Jacobian becomes poorly approximated.

As far as the optimization of adsorption cycles is concerned, three main approaches have been proposed: (i) black-box methods; (ii) simultaneous tailored methods; and (iii) equation oriented methods. Black-box approaches couple a detailed model for the simulation of the adsorption column and an external optimization algorithm, where few independent design and operation variables are optimized. The optimization problem may be extremely challenging because (i) the numerical solvers employed to solve the PDEs of the model may make objective function and constraints noisy and non-smooth, (ii) every cycle simulation needs to reach the CSS condition, making the computational time considerable (about 1-10 min to reach CSS condition). The black-box approach, integrated with meta-heuristic global optimization algorithms, has been adopted by different authors belonging to the adsorption research community. Notably, Boukouvala and Floudas have recently devised a methodology for the single-objective optimization of PSA processes where the problem is formulated as a constrained grey-box problem. The methodology couples in an original approach the use of surrogate functions, deterministic global optimization algorithms (e.g., or depending on the type of problem), and a clustering algorithm to reduce the search space.

In the simultaneous tailored approach, the CSS condition is included as a constraint of the optimization problem in which both design variables and initial states are optimized. In this way, the CSS condition must be guaranteed only at the optimal point with a considerable saving of computational time. On the other hand, since also the initial state variables need to be optimized, the number of optimization variables is larger and a gradient-based optimization algorithm is required. Several approaches are available to deal with the additional CSS constraint: in Kim et al., the difference between the initial and final state of the reactor (temperature, pressure, gaseous component concentration, and the solid component concentration) is constrained to be lower
than a given, very small, positive variable (\(\epsilon\)). On the other hand, other authors\textsuperscript{29} adopt the approach proposed by Ko\textsuperscript{30}, i.e. including \(\epsilon\) as a penalty term in the objective function which is then minimized.

In the equation oriented approach\textsuperscript{14}, the partial differential equations (PDEs) of the column model are expressed as constraints of the optimization problem, such as to impose the CSS condition. Because of the discretization of the PDE equations, the authors have to deal with a large scale nonlinear optimization problem (with up to 50,000 variables and constraints) and adopt an interior point optimization algorithm. However, only simple adsorption models can be tackled with this approach, due to the numerical instabilities introduced by steep front profiles\textsuperscript{31}, the nonlinear and nonconvex nature of the objectives and constraints that may lead to an unfeasible solution or to a local minimum.

As an additional challenge, adsorption processes features multiple contrasting performance. For example, purity and recovery of the target component identify an attainable region of separation, whose upper limit is defined by a Pareto front. A similar behaviour is found for the productivity (defined as the amount of desired product per unit cycle time per amount of adsorbent material), and the energy consumption required for the separation\textsuperscript{32}. While cost evaluation would enable to weigh productivity and energy consumption, it would also introduce significant uncertainties, especially for new materials and processes. Clearly, the identification of the most performing design and operating variables of the process calls for a multiobjective optimization approach. On the other hand, multi-objective optimization is significantly more computational demanding compared to single objective optimization, because of the need of finding a set of optimal solutions.

In this work, we tackle the multiobjective optimization of adsorption cycles. While most of the studies use multiobjective evolutionary algorithms for this purpose\textsuperscript{18,20}, we devise a potentially more efficient algorithm. It is an adaptation of the model-based derivative-free algorithm called Multilevel Coordinate Search (MCS) originally proposed by Huyer and Neumaier\textsuperscript{31}. Notably, the advantage of model-based derivative-free algorithms compared to evolutionary algorithms is the exploitation of local models of the computationally expensive function so as to save evaluations\textsuperscript{34}. In addition, MCS is deterministic, it features convergence guarantees on continuous functions (but only for an infinite number of \textquotedblleft splitting levels\textquotedblright, see Subsection 4.1) and according to the benchmarking work reported in Rios and Sahinidis\textsuperscript{35}, it is the most performing non-commercial algorithm for the global optimization of bound-constrained non-convex smooth problems. In this work, MCS is extended to handle general nonlinear constraints (both algebraic and simulation-based constraints) in addition to bound constraints, and to efficiently identify the set of Pareto optimal solutions (coupling parallel computing with look up tables). We denote this new algorithm as MO-MCS.

This paper is organized as follows: Section 2 describes the multiobjective optimization problem tackled in this work; the model of the adsorption bed and the numerical solution approach are outlined in Section 3; the optimization algorithm is detailed in the Section 4; finally, Sections 5 and 6 show the performance of the MO-MCS algorithm for the optimization of TSA and PSA cycles, respectively.

## 2 Problem statement

The constrained, nonlinear, multiobjective optimization problem tackled in this work can be stated as follows

\[
\begin{align*}
\text{Min } e(\mathbf{x}) \\
\text{Max } Pr(\mathbf{x}) \\
s.t.: \\
\ r_{CO_2} & > r_{CO_2, min} \\
\ \phi_{CO_2} & > \phi_{CO_2, min} \\
\ lb & \leq \mathbf{x} \leq ub \\
\ \mathbf{x} & \in \mathbb{R}^n
\end{align*}
\]

where, \(e\) denotes energy-related objective function, \(Pr\) denotes the productivity (defined as the amount of CO\(_2\) that is produced by a unit mass of sorbent in a unit of time), and \(\mathbf{x}\) is the vector of decision variables (duration of the time steps, volumetric feed flow rate and operating temperatures/pressures); \(r_{CO_2}\) and \(\phi_{CO_2}\) denote respectively the CO\(_2\) recovery (defined as the amount of CO\(_2\) produced over the amount of CO\(_2\) input) and the


CO₂ purity; \( lb \) and \( ub \) define the upper and lower bounds on the decision variables \( x \). Concerning the energy need, the energy consumption is considered when dealing with PSA process, while the exergy consumption is considered for TSA processes, as it combines electric and thermal power consumption; in both cases, the energy need is an indication of the operational costs. On the other hand, the productivity serves as an indirect index of the capital investment of the adsorption processes. Larger productivity enables the use of more compact equipment units for same yield, hence it allows minimizing the equipment cost regardless of the uncertainty of cost related data.

As explained in Section 1, in the black-box optimization approach, the functions \( e(x), Pr(x), r_{CO_2}(x) \) and \( \phi_{CO_2}(x) \) are essentially black-box functions whose output is determined by executing the cycle simulation code (see next section).

3 Modelling cyclic adsorption processes

The cyclic adsorption process is simulated by using an in-house model (named FAST) specifically developed for adsorption processes within fixed beds; the main features of the tool are reported in the following, while details can be found elsewhere \(^{36-38}\). The model consists of a system of partial differential equations (PDE) that account for the mass, energy and momentum balances within the column. The set of PDEs, which is summarized in the supporting information, is solved by discretizing the column along the axial coordinate in finite volumes of constant width. The cycle is described by simulating a single column undergoing each step of the cycle; therefore, the model iterates until cyclic steady state (CSS) conditions are reached, namely once the change in composition, temperature and material balance between two different iterations is below a predefined tolerance, and the overall material balance of the cycle is satisfied. The model is written in FORTRAN, which limits the computation time while also providing the developer with large flexibility. The set of PDEs is solved calling the LSODE solver \(^{39}\), which replaces the ODEs with difference equations and solves them step by step. In our implementation, LSODE uses the backward differentiation formula (BDF) method, which is recommended for stiff problems, and makes use of a modified Newton iteration method to correct the initial estimate of the solution.

The main assumptions of the adsorption model are:

- Radial gradients are negligible
- The gas phase and solid phase are in thermal equilibrium
- Axial dispersion and conductivity are neglected
- Mass and heat transfer coefficients, heat of adsorption and molar heat capacities are assumed to be constant and do not depend on temperature, pressure nor composition
- The mass transfer kinetics are described with the linear driving force (LDF) model.

The model has proven to be a very powerful tool in both reproducing experimental results and designing new adsorption cycles. As far as the experimental validation is concerned, the software was employed for PSA-based separation of H₂/CO₂ mixtures \(^{40}\) and multicomponent H₂/CO₂/N₂ mixtures \(^{41}\) and TSA-based separation of N₂/CO₂ \(^{42}\). Furthermore, the model has been employed for the design of PSA and TSA adsorption cycles \(^{32,37}\) and for the assessment of new sorbents \(^{43,44}\). The column design considered for all calculations matches the column adopted in the P/TSA test rig present at ETH Zurich and thoroughly described elsewhere \(^{41,42}\). It is worth mentioning that when running PSA simulations, the column is assumed to be adiabatic, while a constant heat transfer coefficient between the bed and the column wall is assumed for TSA applications (two constant values are assumed, which mimic stagnant conditions or gas flowing within the bed, respectively). The axial discretization considers volumes of approximately 5 cm width, while the time discretization is implicit and defined by the BDF algorithm. The convergence to CSS condition is guaranteed by ensuring that both the overall mass balances are satisfied and that the variation of temperature and concentration profiles between successive iterations is negligible. The resulting computational time for achieving CSS is between 30 and 180 seconds, for PSA and ranges from 10 to 90 seconds for TSA.
The operational variables that have to be provided as input consist of: (i) time length of each step of the cycle, (ii) gas flow rate (i.e. the gas inlet velocity), (iii) pressure levels in the cycle, especially for PSA, (iv) temperature of the heating and cooling fluid for TSA operation. Even though the simulation tools has proven to be robust in converging, there could be combinations of operational variables that prevent the solver to find a feasible solution. Notably, 2.5% and 5.3% of the simulations did not converge for PSA and TSA, respectively.

4 MO-MCS algorithm

In this section the MO-MCS algorithm is presented. In Subsection 4.1 we briefly review the original MCS algorithm, as proposed by Huyer and Neumaier. In Subsection 4.2 we present the quadratic penalty approach devised to handle nonlinear constraints. In Subsection 4.3, we propose an approach for adapting the original MCS algorithm to multiobjective optimization problems, including an ad-hoc parallel computing architecture tightly coupled with the multiobjective formulation of the problem.

4.1 Multilevel Coordinate Search (MCS) algorithm

Given a bound-constrained problem, formulated as follows:

General bound constrained problem (P1)

\[
\begin{align*}
\text{min } & f(x) \\
\text{s.t. } & lb \leq x \leq ub \\
& x \in \mathbb{R}^n
\end{align*}
\]

Eq.2

where \( x \) denotes the vector of \( n \) decision variables, \( lb \) and \( ub \) are respectively the lower and upper bounds on the variables, and \( f \) is the objective function to be minimized.

The MCS algorithm aims at finding the global minimum of the problem by following a “branching without bounding” approach, which proceeds by dividing the search space into boxes and uses local optimization routines to refine the search in the most promising boxes. Multiple search techniques (i.e., dividing rectangle, coordinate search, “triple search”, model-based trust-region optimization) are iteratively called by the algorithm. For further details, the reader is referred to the work of Huyer and Neumaier.

The MCS algorithm features a convergence guarantee proved by the authors: if the maximum number of splitting levels goes to infinite and the objective function \( f \) is continuous in the neighbourhood of the global minimum, MCS is guaranteed to converge to the global minimum. The basic reason is that the points sampled by the algorithm would form a dense subset of \( \mathbb{R}^n \).

4.2 Extension of MCS to constrained optimization

In order to apply MCS for the optimization of adsorption cycles (see Section 2), it is first necessary to include nonlinear constraints. The general formulation of the resulting constrained problems, as considered in this work, is shown in Eq.3:

Non linear constrained problem (P2)

\[
\begin{align*}
\text{min } & f(x) \\
\text{s.t. } & g_i(x) \leq 0 \quad \forall i = 1: N_c \\
& lb \leq x \leq ub \\
& x \in \mathbb{R}^n
\end{align*}
\]

Eq.3

where \( f \) is the objective function, \( x \) is the set of decision variables, \( g_i \) are the nonlinear constraints and \( N_c \) is the number of constraints.
Several techniques have been developed to convert a nonlinear constrained optimization problem into a series of unconstrained problems, such as penalty methods, barrier methods and augmented Lagrangian methods. We have selected the quadratic penalty strategy because of the following reasons:

- When optimizing an adsorption cycle, the starting solution often does not meet the constraints (desired purity, recovery, etc) and this would affect the effectiveness of barrier methods.
- In order to handle inequality constraints, augmented Lagrangian methods have to convert them into equality constraints by adding slack variables (one slack variable for each nonlinear constraint). These slack variables are added to the list of decision variables causing an increase of the problem dimensionality. While gradient-based methods can easily handle problems with hundreds of variables, the convergence rate of derivative-free methods (like MCS) is significantly affected by the number of variables.
- The nonlinear constraints in the adsorption cycle problem are “relaxable” constraints since some degree of violation can be tolerated. For this reason, the penalty coefficients do not need to be raised to very high values so as to drive to zero the constraints’ violation. Limiting the penalty coefficients allows to moderate the ill conditioning issue of the Hessian matrix of the objective function minimized by MCS. It is worth noting that, although MCS does not directly use the Hessian of the objective function to find the search direction, it builds and optimizes (with an indefinite quadratic programming algorithm) a quadratic model which inherits the ill conditioning of the objective function. For this reason, it is advisable not to adopt excessively high penalty coefficients.

Using the quadratic penalty approach, the constrained nonlinear program (Eq.3) is converted into the bound constrained problem shown in Eq.4.

\[
\text{Bound constrained problem (P3)}
\]

\[
\begin{align*}
\min Q(x, \mu_k) &= f(x) + \mu_k \cdot \sum_{i=1}^{N_c} (\max \left[ g_i(x), 0 \right])^2 \\
\text{subject to} \quad lb &\leq x \leq ub \\
x &\in \mathbb{R}^n
\end{align*}
\]

The original objective function is penalised by adding the summation of the square of the constraint violations times the penalty parameter \( \mu_k \) (k is the iteration index counting the number of times the problem has been solved). Starting from an initial (user-defined) penalty parameter \( \mu_0 \), the bound-constrained problem is solved with the MCS algorithm. Then, the degree of violation of the constraints is checked and, if greater than the desired tolerance, the penalty parameter is increased (\( \mu_{k+1} = \tau \mu_k \) with \( \tau > 1 \)) and the penalized problem is solved again. Convergence of the quadratic penalty approach to a local minimum is guaranteed even if the intermediate k-th penalized problems are not exactly solved (provided that the optimization accuracy progressively decreases tending to zero for k going to infinite). The drawback of this procedure is the need of solving the bound-constrained problem several times with increasing values of the penalty parameter \( \mu_k \).

Due to the high computational cost of the black-box function (i.e., cycle simulation code), each optimization (iteration k) takes several hours making this strategy not practical. On the other hand, to save computational time, one could think of avoiding the intermediate iterations and starting directly with a very high penalty parameter (static penalty approach). Beside the ill conditioning issue of the problem, which may affect the effectiveness of the local search step of MCS, starting directly with high penalty parameters increases the risk of finding local optima or even infeasible solutions, especially if the starting point is infeasible. Our preliminary computational experiments have shown that for some classes of problems with a narrow feasible region it is not only necessary to start with a low \( \mu_0 \) but also increase it very slowly (i.e., adopting \( \tau = 1.15-1.5 \)).

The block-flow diagram of the overall constrained MCS algorithm is shown in Figure 1. In addition to the quadratic penalty approach, the key peculiarities of the algorithm are:

- the use of a look-up table to save evaluations of the black-box function;

\[
\text{Eq.} 4
\]
- a starting point search procedure to select the best starting point (among the ones previously sampled) for each run depending on the value of the penalty parameter $\mu_k$.

Since the simulation of adsorption cyclic processes is computationally intensive, it is essential to save black box function evaluations. Accordingly, we introduced a look up table where the values of objective functions, constraint violations and the corresponding set of decision variables are saved for all function evaluation. When MCS needs to sample a point, it first searches in the look-up table, and, if the desired point has not been yet simulated, it executes the adsorption cycle simulation code. The adoption of look-up tables within MCS is particularly advantageous because:

1) If the increase of $\mu$ from iteration $k$ to the next, $k+1$, is not excessive, the global minimum of the penalized problem will likely remain in the same region.

2) If the same initialization list of splitting points is passed to MCS, in the first Sweep iteration it will proceed by dividing the feasible region in the same way, sampling the same base points as in the previous ($k-1$) optimization run. This occurs because MCS is a deterministic algorithm without random components.

Along with look-up tables, a procedure was designed in order to select the best starting point for the subsequent optimization run ($k+1$) with increased penalty parameter. When an instance of problem (P2) is solved with fixed penalty $\mu_k$, the algorithm may return an intermediate suboptimal or even infeasible solution (i.e., which does not meet all the constraints $g_i$ because the penalty parameter is too low). When the penalty parameter is updated at step $k+1$ ($\mu_{k+1}$), the best feasible point in the look-up table is selected as starting point for the $k+1$ iteration.

4.3 Extension of MCS to multiobjective optimization and parallelization strategy

A widely-used technique for generating a Pareto set in multiobjective optimization is the weighted sums (WS) approach. The basic idea is to convert the multiobjective problem into a sequence of single objective problems, where a convenient convex combination of the two objective functions is minimized. However, the WS method has two main drawbacks: (i) it cannot find non-convex Pareto frontiers, (ii) obtaining an even distribution of solutions on the Pareto frontier is difficult. Ko proposed an extension of the WS method, which is able to return well-spaced Pareto points, but to the cost of several function evaluations, with therefore scarce applicability to computationally intensive problems. Another technique for generating the Pareto set is the Normal-Boundary Intersection (NBI) method. On the one hand, the NBI method can achieve fairly

---

Figure 1 Block-flow diagram of the overall constrained MCS algorithm
uniform distribution of Pareto-optimal solutions and can treat problems with non-convex Pareto fronts by imposing equality constraints along equally spaced lines. On the other hand, including equality constraints makes the optimization problems more challenging and the NBI approach may return dominated solutions (not-Pareto optimal)\(^4^9\). These two issues have been addressed by the Normalized Normal Constraint (NNC) method\(^5^0\). The NNC method converts the multiobjective problem into a sequence of single objective nonlinear constrained problems (NLP) subject to normalized normal inequality constraints which limit the feasible objective space. As shown in Figure 2B, the \(j\)-th Pareto point is obtained by minimizing one of the two objectives subject to the \(j\)-th normalized normal constraint. Thus, a set of equally spaced Pareto points is obtained by performing several single objective optimizations with shifted normalized normal constraints. Compared to other approaches, NNC can deal with non-convex Pareto fronts and return equally spaced Pareto solutions without the risk of finding dominated solutions\(^5^0\).

Owing to its above-mentioned advantages, we selected the NNC approach to convert the constrained bi-objective optimization problem (Eq.5) into a set of single objective constrained problems which can be efficiently tackled with the MCS algorithm. Moreover, we devised an implementation capable of exploiting multiple processors (workers) in parallel computing while maximizing the benefit of the look-up table.

Given the constrained nonlinear bi-objective in Eq. 5:

Constrained nonlinear bi-objective problem (P4)

\[
\text{min}\{f_1(\mathbf{x}), f_2(\mathbf{x})\}
\]

\[
\text{s.t.}\ g_i(\mathbf{x}) \geq 0 \quad i = 1: N_C
\]

\[
\text{lb} \leq \mathbf{x} \leq \text{ub}
\]

\[
\mathbf{x} \in \mathbb{R}^n,
\]

Eq. 5

\[
\text{Figure 2 implementation of the NNC method represented in the objective space (o.s.)}
\]

Our implementation of the NNC method consists of the following steps:

1. **Anchor points calculation**. The extreme points of the Pareto front (anchor points) are found solving two single objective instances (A1 and A2) of problem P4, minimizing respectively one of the two objectives while neglecting the other.

Problem (P5): Minimization of \(f_1\)

\[
\text{min} \ f_1(\mathbf{x})
\]

\[
\text{s.t.}\ g_i(\mathbf{x}) \geq 0 \quad i = 1: N_C
\]

\[
\text{lb} \leq \mathbf{x} \leq \text{ub}
\]

Eq. 6
\( x \in \mathbb{R}^n \)

Problem (P6): Minimization of \( f_2 \)
\[
\min f_2 (x)
\]
s.t.
\[
g_i(x) \geq 0 \quad i = 1: N_c
\]
\[
lb \leq x \leq ub
\]
\( x \in \mathbb{R}^n \)

Let us define \( f_{A1} = [f_1(x_{A1}), f_2(x_{A1})] \) and \( f_{A2} = [f_1(x_{A2}), f_2(x_{A2})] \) the corresponding optimal solutions obtained. The line joining them is called “utopia line” (see Figure 2 A).

2. **Normalization.** The values of the objective functions are normalized according to the utopia points:
\[
l_1 = f_1(x_{A2}) - f_1(x_{A1}) \quad \text{Eq.8}
\]
\[
l_2 = f_2(x_{A1}) - f_2(x_{A2}) \quad \text{Eq.9}
\]
\[
\hat{f}(x) = \left(\frac{f_1(x) - f_1(x_{A1})}{l_1}, \frac{f_2(x) - f_2(x_{A2})}{l_2}\right) \quad \text{Eq.10}
\]

3. **Utopia line vector definition.** The Utopia line vector \( N_1 \) is a vector connecting the two extreme points of the Pareto front, defined in Eq.11.
\[
N_1 = f_{A2} - f_{A1} \quad \text{Eq.11}
\]

4. **Definition of the number of Pareto-points.** From this step on, it is necessary to set the desired number of points on the optimal Pareto-front, in addition to the two extreme points. According to the proposed approach, every Pareto point is the solution of a single objective optimization problem, and each problem is assigned to a specific worker. In order to maximize the benefit of parallel computing, we recommend to select a number of internal Pareto points multiple of the number of workers of the machine (\( N_W \)). In this way, the calculation of the same number of internal Pareto points (here called “number of worker turns”, \( N_{WT} \)) is assigned to each worker. Then, the total number of Pareto points (\( N_{PP} \)) (included the two extreme points) and the normalized increment \( \delta \) can be obtained as shown in equation Eq.12 and 13.
\[
N_{PP} = (N_W \cdot N_{WT}) + 2 \quad \text{Eq.12}
\]
\[
\delta = \frac{1}{N_{PP} + 1} \quad \text{Eq.13}
\]

5. **Utopia line points generation.** Utopia line points \( (X_j) \) are a series of equally spaced points along the utopia line.
\[
X_j = \alpha_{1,j} \cdot f_{1}^* + \alpha_{2,j} \cdot f_{2}^* \quad \text{Eq.14}
\]
\[
0 \leq \alpha_{1,j} \leq 1 \quad \text{Eq.15}
\]
\[
\alpha_{1,0} = 0 \quad \text{Eq.16}
\]
\[
\alpha_{1,j+1} = \alpha_{1,j} + \delta \quad \forall j = 1: N_{PP} - 1 \quad \text{Eq.17}
\]
\[
\alpha_{2,j} = 1 - \alpha_{1,j} \quad \forall j = 0: N_{PP} \quad \text{Eq.18}
\]

We note that \( \alpha_{1,j} \) is incremented by \( \delta \) between 0 and 1. Values \( \alpha_{1,0} = 0 \) and \( \alpha_{1,N_{PP}} = 1 \) are discarded, since they would lead to the same Pareto points found in step 1.

6. **Generation of the Pareto points**
Using the set of evenly distributed Utopia line points, a corresponding set of Pareto points is generated by solving a succession of optimization runs. J-th optimization run yields a Pareto optimal point and is associated to utopia line point \( X_j \) by the addition to the original problem P4 of a constraint normal to the utopia vector and passing through the j-th utopia point. In our approach, each optimization run is assigned to a specific worker and two different types of NNCs (\( h^4 \) and \( h^2 \)) are considered simultaneously as shown in Figure 2C. Each NNC type starts from an opposite end of the Pareto front and cuts off opposite sides of the objective feasible space. Every optimization run is subject to either \( h^4 \)or \( h^2 \) NNC. Since each optimization run is carried out by one and only one worker, we can distinguish between

- \( N_{pp}/2 - 1 \) optimizations runs, which are associated to the \( N_{pp}/2 \) utopia line points closer to \( A_1 \) and are subject to \( h^2 \) constraint while minimizing \( f_2 \), represented by problem P7 (Eq.19)
- The other \( N_{pp}/2 - 1 \) points, closer to \( A_2 \) which will be subject to \( h^4 \) constraint while minimizing \( f_1 \), represented by problem P8 (Eq.20).

\[
\text{Problem (P7): } \forall \ j: 1 \leq j \leq N_{pp}/2
\]
\[
\begin{align*}
\min & \ f_2(x) \\
\text{s.t.} & \quad g_i(x) \geq 0 \quad i = 1:N_c \\
& \quad h_j^2(x) = -N_j \cdot (f(x) - X_j)^T \geq 0 \\
& \quad lb \leq x \leq ub
\end{align*}
\]

\[
\text{Problem (P8): } \forall \ j: N_{pp}/2 < j \leq N_{pp} - 1)
\]
\[
\begin{align*}
\min & \ f_1(x) \\
\text{s.t.} & \quad g_i(x) \geq 0 \quad i = 1:N_c \\
& \quad h_j^4(x) = N_j \cdot (f(x) - X_j)^T \geq 0 \\
& \quad lb \leq x \leq ub
\end{align*}
\]

As for the additional NNC constraint added in each optimization run, MO-MCS adopts the same methodology used for the black box constraints, and outlined in section 4.2.

Clearly, this overall approach makes full use of the algorithm parallelization: half of the workers are assigned problems of type P7, and the other half problems of type P8. As illustrated in Figure 3, the optimization runs tackled in the first worker turn lead to the Pareto points closer to the anchor points. Anchor point A1 is used as starting point for problems of type P7, while anchor point A2 is the starting point for problems of type P8. We observe that this technique pledges the feasibility of the starting point for every optimization run, for both problems of type P7 and P8. Moreover, for a fixed number of workers, the proposed starting point of each optimization run allows for a faster identification of the next Pareto point, especially when compared to single NNC strategy. If more worker turns are needed, the starting points for both classes of problems are updated, choosing the last Pareto points derived in the previous worker turn. In this way, the Pareto front is built symmetrically from the ends towards the centre, as represented in Figure 3. The complete structure of the algorithm is represented in Figure 4. A look up table is included, where information about already performed black box function evaluation are stored. Whenever a new function evaluation is needed, the algorithm first scans the look up table to check if the point has already been explored by worker in the same worker turn or in preceding worker turns. If so, results are retrieved from the look up table and the computational time for a black box function evaluation is saved. The last worker turns operate in more beneficial state, since a higher number of points has been explored and is available in the look up table.
Figure 3: Example of the generation sequence of the Pareto front for a case with 4 workers and 2 worker sweeps. Each shape corresponds to a specific worker. The anchor points are derived in parallel by two different workers and marked in red. The points derived in parallel in the first worker sweep are marked in blue and the ones derived in parallel during the second worker sweep are marked in green. Dashed lines link the starting point used to find each Pareto point. Points originated from the same anchor point are subject to a NNC constraint with the same direction.
Figure 4 General scheme for the parallelization of MO-MCS. After the initialization step, the two extreme Pareto points are derived in parallel. The problem is normalized, and the utopia line points are generated. The starting points for the first worker turn are assigned. Index \( j \) spans the instances of the problem to be solved, each assigned to a specific worker and associated to a specific NNC. In each worker turn \( z \), the first \( N_w/2 \) instances tackled are subject to \( h^1 \) while minimizing \( f_1 \), while the other \( N_w/2 \) problems are subject to \( h^2 \) while minimizing \( f_2 \). After each worker turn, the starting points are updated and a new worker turn is performed, until \( z = N_{WT} \). During the search for each Pareto point (marked with *), each worker can read information from the look-up table (which is common to all workers) and write the results of new sampled points. In this way, each worker can access information of the black box function evaluations performed by other workers in the same work turn, as well as the ones performed in previous worker turns.
As far as the algorithm parameters of MO-MCS are concerned, good results were obtained with those reported in Table 1 (and adopted for all the test cases of this work). It is worth noting that the number of worker turns is not critical for the quality of the results, and it is set depending on the desired number of Pareto points on the front. It is also worth stressing that a large number of additional Pareto points is found during the execution of the algorithm, which can be retrieved from the look-up tables without any additional computational expense. However, no guarantee on the feasibility, number and spacing of such additional points exist. This means that it is likely, but not sure, to obtain a finer approximation of the Pareto curve than the number of points set as input.

**Table 1 Suggested parameter tuning for MO-MCS**

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Suggested value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Initial value of the penalty parameter ((\mu_0)) for extreme Pareto points search:</td>
<td>5</td>
</tr>
<tr>
<td>Initial value of the penalty parameter ((\mu_0)) for internal Pareto points search:</td>
<td>0.1</td>
</tr>
<tr>
<td>Initial value of the penalty parameter ((\mu_{0,NCC})) for NNC constraint:</td>
<td>0.1</td>
</tr>
<tr>
<td>Penalty parameters increase rate ((\mu_{k+1}/\mu_k))</td>
<td>1.5</td>
</tr>
</tbody>
</table>

It is worth noticing that the search for extreme Pareto points involve one less constraint (the NNC) and hence is less difficult, so the initial value of the penalty factor can be higher than that used for the calculation of the internal Pareto points (as suggested in Table 1). When possible, normalization of the constraints is recommended so as to balance their weights.

If the user adopts a poor normalization of the constraints or too high/low initial values of the penalty parameters, one may obtain:

- a perfectly spaced but sub-optimal Pareto curve, if the weight of the NNC constraint is too large compared to the weight of the other constraints of the problem.
- a poorly spaced Pareto curve, if the weight of the NNC constraint is too small compared to the one of the other constraints of the problem.

5 MO-MCS benchmarking

In this section, the performances of MO-MCS are compared against other well-known, derivative-free commercial algorithms. A similar benchmarking procedure, yet for an earlier version of MO-MCS, is reported in Joss et al. \(^51\). Because of the novel features, which improve considerably the computational performances of the algorithm, an updated benchmarking is reported here. Moreover, in this work, we have modified (e.g. enabling parallel computing) and tuned all considered algorithms to the best of our capabilities.

The test problem is a five-step TSA cycle, as considered in Joss et al. \(^51\), where constraints on the purity and recovery of the separated CO\(_2\) are considered. The objective functions are the productivity of the cycle (to be maximized), and the energy consumption (to be minimized). The decision variables are: the duration of four out of five time steps (two time steps require equal duration). As far as the benchmarking procedure is concerned, the process details of the test problem are not relevant, and the reader is referred to \(^51\) for further details.

The algorithms selected for the benchmarking are:
1. Non Sorting Genetic Algorithm II (NSGA II), a genetic algorithm presented by Deb et al. in 16, which can handle multiobjective constrained problems with a fixed penalty approach.

2. Multi-Objective Particle Swarm Optimization (MOPSO), proposed by Coello at al. in 52. The original algorithm presents no built-in parallel computing approach. For the sake of this article, MOPSO algorithm was updated to include such skill. In particular, the evaluation of each particle in the current population is parallelized, for both the initialisation step and the subsequent repository updates.

3. Direct Multi Search (DMS), presented by Custodio et al. in 53.

Both MOPSO and DMS algorithms have no standard routine to deal with constrained problems, so the same fixed penalty strategy available in NSGA II is extended to the other benchmarking algorithms. The test problem is first solved with the MO-MCS algorithm, then, the number of function evaluations required by the MO-MCS (reported in Table 2) is set as stopping criterion for the other algorithms. Such upper limit is obtained by setting broad stopping criteria for the MO-MCS (i.e. max 50 function evaluations for each fixed penalty run and max 150 function evaluations for each Pareto point), and observing the algorithm convergence within these limits. The fixed penalty approach adopted to make NSGA II, MOPSO and DMS suitable for constrained problems calls for a fine-tuning of the penalty parameter: too low a value may lead to unfeasible results, while too high a value may ill-condition the problem and lead to suboptimal results. In order to perform a fair comparison, several values of penalty factor have to be tested, since it is not possible to determine a priori an effective value. The values considered for the penalty parameter are: 0, 2, 5, 20, 250. As expected, lower values of the penalty parameters perform better, but runs with penalty factor lower than 5 are discarded, because they lead to unfeasible results. We report here only the best results obtained. Notably, MO-MCS features an auto-adaptive strategy for the penalty factor to solve such tuning issue, that may perform well enough without tuning. While we don’t report here the time for tuning, as it is very much dependant on the user, it is useful to keep this mind when tackling a similar optimization problem.

Since the NSGA II and MOPSO are stochastic meta-heuristic algorithms, five runs of each algorithm were performed and the best three runs are reported in Figure 5 A and B. On the other hand, multiple runs are not necessary for DMS as it is a deterministic algorithm. The run with the best penalty parameter is reported in Figure 6. All the results are plotted along with MO-MCS outcome. All the Pareto fronts are plotted in dimensionless form with the normalized productivity on the x-axis and the normalized energy consumption on the y-axis.

### Table 2 Benchmarking optimization parameters.

<table>
<thead>
<tr>
<th></th>
<th>MO-MCS</th>
<th>NSGAII</th>
<th>MOPSO</th>
<th>DMS</th>
</tr>
</thead>
<tbody>
<tr>
<td>Workstation used to perform optimization runs</td>
<td>Intel ® Xeon ® CPU 2.79 GHz</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Maximum number of function evaluations allowed (stopping criterion)</td>
<td>1950</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Number of workers</td>
<td>4</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Number of well-spaced Pareto points</td>
<td>10</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Population size</td>
<td>50</td>
<td>50</td>
<td>100</td>
<td></td>
</tr>
<tr>
<td>Parallel computing</td>
<td>YES, available</td>
<td>YES, adapted</td>
<td>NO</td>
<td></td>
</tr>
</tbody>
</table>

14
The results reported in Figure 5 A show that MO-MCS outperforms NSGAII, since with a single run it can provide a Pareto front which is quite well spaced, wider and it dominates all the Pareto fronts obtained with five runs of NSGAII, i.e., for a fixed productivity MO-MCS finds points with lower (or equal) energy consumption. The computational time required for a single run of NSGAII is 24% shorter than the time required from MO-MCS (see Table 3), thanks to the more efficient use of parallel computing: in fact, MO-MCS includes a first step in which just two workers are exploited. On the other hand, it is important to remark that NSGAII required multiple runs to obtain a good approximation of the Pareto front, thus its total required computational time is considerably higher than a single run of MO-MCS.

In Figure 5 B, the same Pareto front obtained from MO-MCS is compared to the results obtained from MOPSO. Three different runs are reported for MOPSO, and in this case, the best runs are obtained adopting the same penalty factor. Even if MOPSO is able to find some points laying on the optimal Pareto front, MO-MCS provides a better distributed Pareto front with larger extension and number of Pareto points. The computational time required by a single run of MOPSO is slightly shorter than the one of NSGA II, and indeed shorter than the one for MO-MCS, as shown in Table 3. However, due to the need of making multiple runs of MOPSO, MO-MCS is much more efficient from a computational point of view.
Figure 6 Normalized Pareto fronts for DMS algorithm compared to the normalized Pareto front obtained with MO-MCS. Both the output from MO-MCS (filled dots) and the additional Pareto points retrieved from the look up tables (empty dots) are considered. In addition, the initial guess adopted for all runs is reported. The shape of the Pareto front is because higher productivities are associated with higher energy consumptions, as expected.

In Figure 6, the Pareto fronts obtained from MO-MCS and DMS are compared. DMS can find some points laying on the optimal Pareto front identified by MO-MCS, although its Pareto front is not as well spaced and extended. From the computational time perspective, DMS is the slower among all the considered optimizer because the code used for the tests does not exploit parallel computing (although Custodio et al.\textsuperscript{53} suggest the possibility of parallelizing the poll steps, this option was not available in the most updated version (0.3) of the code). Unfortunately, even if the computation time is drastically reduced, we found that in such case the resulting Pareto front is completely dominated by the one obtained with sequential computing. This is due to the fact that in the sequential computing approach it is possible to determine the non-dominated Pareto set after every single function evaluation (i.e. discard the dominated particles), while in the parallel computing approach such dominance relationships can be established only after each polling step, reducing the amount of information available during the polling step.

The performances of the four algorithms are quantified according to Zitzler’s indices\textsuperscript{55}, and are reported in Table 3. Index M1 expresses the average distance from the Pareto optimal set, which is approximated with the envelope of all the Pareto fronts derived in the benchmarking procedure. Index M3 use the maximum extent in each dimension to measure the spreading of the Pareto front. Index M2 aims at measuring the number of well spaced points in a Pareto front. This index includes a parameter $\sigma$, which is the minimum spacing in order to consider a pareto to be well spaced. In this paper we assume such spacing according to Eq.21, i.e. the minimum satisfactory spacing is the one between $N_{\text{particles}}$ equidistributed along a line connecting Pareto extreme points in the normalized plane.

$$
\sigma = \frac{\sqrt{2}}{(N_{\text{particles}}-1)} \tag{Eq.21}
$$

Since for NSGA-II and MOPSO the performance of the Pareto front will vary for the different runs, we decided to evaluate the performance indices based on the outer envelope of all Pareto points considering the three different optimization runs.
Table 3: Benchmarking optimization results. Index M1 measure the distance of the obtained Pareto front from the overall best envelope considering all Pareto fronts. Index M2 expresses the number of well-spaced particles in the Pareto fronts, while index M3 measures the extension of the Pareto fronts.

<table>
<thead>
<tr>
<th></th>
<th>Index M1, %</th>
<th>Index M2</th>
<th>Index M3</th>
<th>Average computational time for a single run, h</th>
</tr>
</thead>
<tbody>
<tr>
<td>MO-MCS</td>
<td>0.19%</td>
<td>183</td>
<td>1.41</td>
<td>1.81</td>
</tr>
<tr>
<td>NSGA II (best envelope on 5 runs)</td>
<td>2.59%</td>
<td>28</td>
<td>1.18</td>
<td>1.37</td>
</tr>
<tr>
<td>MOPSO (best envelope on 5 runs)</td>
<td>1.21%</td>
<td>24</td>
<td>1.26</td>
<td><strong>1.33</strong></td>
</tr>
<tr>
<td>DMS</td>
<td>1.87%</td>
<td>174</td>
<td>1.33</td>
<td>3.02</td>
</tr>
</tbody>
</table>

Even if the competing algorithms are advantaged by the larger number of runs executed for each instance, MO-MCS performs better in terms of Pareto optimality (M1), spacing (M2), and width (M3). MOPSO ranks second in optimality of the returned Pareto curve (second lowest M1), but its spacing and spread are poorer than the ones achieved with MO-MCS. DMS trails MO-MCS according to M3 and M2, thanks to the large number of non-dominated points, however it shows longer computational time and worse M1 than MOPSO. NSGA-II returns Pareto curves with the lowest extension, and hence ranks last according to index M3. MO-MCS also outperforms NSGA-II according to M1, since also the distance from the best Pareto envelope is penalised by the reduced extension of the Pareto curve.

6 Case studies

In the following, we apply MO-MCS to more challenging case studies, where fixed bed cycles are designed to fit the purpose of CCS. The considered processes are: (i) a TSA cycle for the separation of CO$_2$ from ambient pressure fluegas, and (ii) a PSA cycle for the separation of H$_2$ and CO$_2$ from a pressurized syngas.

Not only do these two processes present key features that call for rigorous yet challenging optimization, but they are also of great interest as alternative separation technologies to the more consolidated liquid scrubbing.

Notably, the two case studies allow to test the MO-MCS with two processes that, while sharing features of adsorption-based technologies, significantly differ in terms of driving force and typical time length of the cycle steps. Indeed, a TSA cycle is controlled by heat exchange, which may take several minutes to bring the reactor to the desired conditions, while a PSA cycle is controlled by pressure change, which is notably fast. This also affects the resolution of the cycle simulation, as the TSA is inherently longer, while the sharp change of fronts in PSA makes the resolution more stiff.

6.1 TSA cycle

The considered TSA cycle was presented in Joss et al. (referred to as Cycle D) as an advanced cycle for the separation of CO$_2$ starting from a binary CO$_2$-N$_2$ mixture at ambient pressure. The cycle is shown in Figure 7. Compared to the simplest TSA cycle, which comprises an adsorption, a heating and a cooling step, this cycle possesses the following additional steps: (i) a rinse step, which is conducted after the adsorption by recycling a CO$_2$-rich gas, (ii) a preliminary heating step following the rinse, where the effluent stream is recycled to the adsorption step, and (iii) a purge step, which is carried out after the heating by purging the column with N$_2$, and which provides the gas to the rinse. Joss et al. showed that this cycle provides excellent separation performance with energy and productivity similar to amine scrubbing. It is worth stressing that with this work we don’t aim at investigating the TSA process comprehensively (e.g. no drying or scheduling has been considered), but rather at explaining the physical rationale behind the optimal points identified by MO-MCS.
The fluegas entering the column has the following molar composition:

- N₂: 88%
- CO₂: 12%

The considered sorbent material is as in Joss et al. 32, i.e. a commercial zeolite 13X. Heating and cooling are carried out with an external thermal fluid so as to preserve the CO₂ purity. Eight decision variables are considered in the optimization: the time steps of adsorption, heating1, heating2, rinse – which constraint also the purge step – and cooling, the temperature of the heating fluid, the mass flow rate of inlet fluegas per column, and the mass flow rate of nitrogen used as purge gas. The list of variables along with the corresponding range of values, is reported in Table 4.

Table 4: List of the decision variables for the optimization of the TSA cycle and their respective bounds.

<table>
<thead>
<tr>
<th>Decision variables</th>
<th>Upper – Lower bounds values</th>
<th>Initial guess</th>
</tr>
</thead>
<tbody>
<tr>
<td>Time steps for:</td>
<td></td>
<td></td>
</tr>
<tr>
<td>1. Adsorption, ( t_{ads} ) [s]</td>
<td>1. 10</td>
<td>1. 4'000</td>
</tr>
<tr>
<td>2. Recycle (and purge, as they are equal), ( t_{rec} ) [s]</td>
<td>2. 5</td>
<td>2. 1'000</td>
</tr>
<tr>
<td>3. Heating 1, ( t_{H1} ) [s]</td>
<td>3. 1</td>
<td>3. 2'000</td>
</tr>
<tr>
<td>4. Heating 2, ( t_{H2} ) [s]</td>
<td>4. 20</td>
<td>4. 4'000</td>
</tr>
<tr>
<td>5. Cooling, ( t_{C} ) [s]</td>
<td>5. 20</td>
<td>5. 4'000</td>
</tr>
<tr>
<td>6. Temperature of the heating fluid, ( T_{H} ) [K]</td>
<td>6. 380</td>
<td>6. 450</td>
</tr>
<tr>
<td>7. Flow rate of inlet fluegas, ( m_{FG} ) [m³/s]</td>
<td>7. 3.5×10⁻⁴*0.7</td>
<td>7. 3.5×10⁻⁴*1.3</td>
</tr>
<tr>
<td>8. Flow rate of the N₂ recycled to purge, ( m_{N2} ) [m³/s]</td>
<td>8. 3.5×10⁻⁴*1.3</td>
<td>8. 3.5×10⁻⁴*1.3</td>
</tr>
</tbody>
</table>

Problem formulation
The optimization problem for the TSA test case is formulated in Eq.22. The objective functions considered are the exergy equivalent (\( e^{ex}_{TSA} \)) of the cycle thermal power demand \( Q_{TSA} \) (specific to the kg of CO₂ separated), as defined in Eq.23, and the productivity. \( e^{ex}_{TSA} \) is selected to take into account the different availability content of heat at different temperature.
In term of separation performance, the TSA optimization is constrained to achieve a CO$_2$ purity ($\phi_{CO_2\_TSA}^{\min}$) of 96% and a CO$_2$ recovery ($r_{CO_2\_TSA}^{\min}$) of 90%.

$$\begin{align*}
\begin{cases}
\text{Min } ex^{TSA}(x) \\
\text{Max } Pr^{TSA}(x)
\end{cases}
\end{align*}$$
Subject to:

$$
\begin{align*}
& r_{CO_2} > r_{CO_2\_TSA}^{\min} \\
& \phi_{CO_2} > \phi_{CO_2\_TSA}^{\min} \\
& lb \leq x \leq ub \\
& x \in \mathbb{R}^n
\end{align*}$$

Eq.22

$$ex^{TSA} = \dot{Q}_{TSA} \cdot \left(1 - \frac{T_0}{T_H}\right)$$

Eq.23

The parameters adopted to run the MO-MCS algorithm for the TSA case study are the same suggested in section 5, except for the ones reported in Table 5: due to the higher complexity of the case study compared to the benchmarking problem, the maximum number of function evaluations is increased both for the fixed penalty runs and overall. Besides, a finer mesh for the Pareto front is specified increasing the number of optimal Pareto points. The optimization run is carried out with the same hardware used for the benchmarking of MO-MCS (see Table 5).

**Table 5: Mo-MCS parameters for the TSA study case**

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Adopted value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of workers</td>
<td>8</td>
</tr>
<tr>
<td>Number of worker turn</td>
<td>2</td>
</tr>
<tr>
<td>Number of Pareto points</td>
<td>18</td>
</tr>
<tr>
<td>Maximum number of function evaluation for a fixed penalty parameter ($N_{max_k}$)</td>
<td>150</td>
</tr>
<tr>
<td>Maximum number of function evaluation for a single Pareto point search ($N_{max_tot}$)</td>
<td>800</td>
</tr>
</tbody>
</table>

**Results**

The overall computational time required to run the TSA optimization is 5.4 hours. The total number of function evaluations required is 10,914. On average, each Pareto point search requires 606 function evaluations. The optimization run leading to the minimum exergy cost configuration is the most challenging, and requires 1,067 function evaluations, which correspond to 6 fixed penalty iterations, meaning it is more difficult to find feasible solutions close to the minimum of exergy cost for the selected initial guess. When considering all the Pareto points, approximately 12% of the objective values are retrieved from the look up tables. However, no clear trend can be established between the number of function evaluations, the percentage of values retrieved from the look up tables and the position of the point along the Pareto curve. For the vast majority of the Pareto points, the resulting number of function evaluations is considerably lower than the maximum allowed, meaning that MO-MCS is able to find feasible optimal solutions within the stopping criterion adopted.

In Figure 8, the Pareto curve resulting from the optimization is represented. Remarkably, all points satisfy purity and recovery constraints. The obtained Pareto curve is smooth, the spacing of the Pareto points is very...
satisfactory, and a clear trade-off between exergy cost and productivity can be easily identified. The productivity ranges from $23.1 \text{ kg CO}_2/(\text{ kg bed h})$ to $106.9 \text{ kg CO}_2/(\text{ kg bed h})$, while the energy cost varies between $1.18 \text{ MJ/(kg CO}_2)$ and $1.49 \text{ MJ/(kg CO}_2)$. The objective value is reported also for the selected starting point, which is suboptimal and not close to any Pareto point.

Figure 8: Pareto curve resulting from the MO-MCS run. The filled dots are the points returned directly from the algorithm, while the empty circles are the non-dominated points found in the look-up tables after MO-MCS is executed.

It is worth noting that the black-box model optimization approach allows to evaluate how the decision variables affect the performance. Indeed, this can be done by considering the value of the decision variables for the Pareto optimal points. Such results are shown in Figure 9. First, the variable trends show that there are multiple, not trivial, combinations of the decision variables that are optimal for the system. These combinations would be hardly identified without a proper optimizer. While some variables show a clear trend moving along the pareto (from minimum to maximum productivity), others do not. In the latter cases it might be possible to identify a range of optimal values. As shown in Figure 9 B, the duration of the heating 2 step is maximum in the low exergy consumption region, while it decreases to achieve higher productivity. Indeed, longer heating times allow the bed to reach a higher temperature, hence a larger amount of desorbed CO$_2$. In parallel, the hot level temperature (Figure 9 F) is maximum in the high productivity region and decreases for low exergy consumptions. This because a higher temperature level implies a higher exergy cost of the thermal power required during heating steps while leading to faster heating. Although no clear trend emerges for the duration of the adsorption and the recycle (purge) step, (Figure 9 E) the optimizer identifies a narrow optimal range
The trends for the feed and N\textsubscript{2} purge flow rate are not reported in Figure 9 since they are constant for all optimal points, hitting the lower and upper bound of the considered optimization range, respectively. Additional information about the influence of the constraints on the Pareto curve are reported with a colour code in Figure 9. The constraint on the purity of the CO\textsubscript{2} stream produced is active along the whole Pareto curve, meaning that it is limiting the performance of the considered adsorption cycle. On the other hand, the constraint on the CO\textsubscript{2} recovery is activated without any clear trend. It is worth noting that the optimizer requires more function evaluations for those points where the constraint on the CO\textsubscript{2} recovery is active.

Figure 9 Trends of the decision variables for the TSA case study. The trend of each decision variable is plotted as function of the normalized productivity. A: Time step heating 1, B: Time step heating, C: Time step purge, equal to time step recycle, D: Time step cooling, E: Time step adsorption, F: Hot temperature level. The empty circles represent optimal solutions were no constraint is active, while filled circles represent points where the constraint on CO\textsubscript{2} recovery is active, green smaller circles represent points where the constraint on CO\textsubscript{2} purity is active.

Scaling of parallel computing performance
In order to understand the effect of the number of workers used to run MO-MCS on the performance of the algorithm itself, a sensitivity analysis is performed: the same TSA case study is solved employing 2, 4, 8, and 16 workers. The number of Pareto points is kept fixed and equal to 18, so the number of worker turns is adjusted consequentially. The starting point and other algorithm parameters are not varied with respect to the previously presented TSA case study. The main results of the sensitivity analysis are reported in Table 6. While the overall number of function evaluation is almost constant, it is possible to identify a clear trend for the computational time, which decreases when increasing the number of workers. The case where all the internal Pareto points are calculated simultaneously (case using 16 workers) leads to the minimum computational time. However, the relative advantage of increasing the number of workers is progressively reduced according to the well-known Amdahl’s Law\textsuperscript{57}. The reasons limiting the speedup are (i) the calculation of the two anchor points can exploit only two workers (thus its computational time is constant when increasing the number of workers), (ii) the serial code sections which cannot be parallelized, (iii) communication overheads between processors, (iv) load imbalance between processors. If more than 16 workers are available, no speedup is expected and it is advantageous to increase the number of calculated Pareto points.
Concerning the use of look-up table, between 12-21% of the function evaluations are retrieved with a considerable saving of computational time. The minimum use of the look-up table occurs in the run with 8 workers because of the following trade-off:

- for larger number of workers (e.g. 16), several workers start from the same starting point and hence tend to follow the same search path, at least in the initial phase of the optimization. When the same points are sampled by different workers (not simultaneously), the look-up table is efficiently used.
- for smaller number of workers (e.g. 2, 4), the successive worker turns benefit from the points sampled in the previous worker turns.

<table>
<thead>
<tr>
<th>Number of parallel workers</th>
<th>2</th>
<th>4</th>
<th>8</th>
<th>16</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of Pareto points</td>
<td>18</td>
<td>18</td>
<td>18</td>
<td>18</td>
</tr>
<tr>
<td>Number of Worker turns</td>
<td>8</td>
<td>4</td>
<td>2</td>
<td>1</td>
</tr>
<tr>
<td>Total number of points sampled</td>
<td>11,064</td>
<td>11,194</td>
<td>10,914</td>
<td>11,064</td>
</tr>
<tr>
<td>Fraction of points retrieved from the look up table</td>
<td>19%</td>
<td>21%</td>
<td>12%</td>
<td>17%</td>
</tr>
<tr>
<td>Computation time, h</td>
<td>13.64</td>
<td>9.03</td>
<td>5.41</td>
<td>4.92</td>
</tr>
<tr>
<td>Computational time reduction</td>
<td>-</td>
<td>-34%</td>
<td>-58%</td>
<td>-64%</td>
</tr>
</tbody>
</table>

6.2 PSA cycle

The PSA cycle was synthesized specifically for this study, with the aim of providing both the light and heavy product (i.e. H₂ and CO₂, respectively) at the purity of interest. Accordingly, the complexity of the cycle increases: following the adsorption step and three pressure-equalizations, which allow to increase the recovery, two blowdown steps (BD) are adopted to regenerate the sorbent. The first blowdown provides a waste product that contains most of the impurities and a remarkable amount of CO₂. The second blowdown provides the CO₂ product. The backpressure regulator is fixed at 1.5 bar for the first blowdown step, 1 bar for the second blowdown step. The actual pressure reached inside the reactors may be higher and is determined by the duration of the respective blowdown step, and is hence an output of the model. The resulting cycle is shown in Figure 10. The increased complexity of the adsorption process, along with the rapid dynamics of pressure swing cycles, make the optimization particularly challenging, yet also particularly suitable for testing the algorithm. In this case the decision variables, which are reported in Table 7, are: the time steps of adsorption, BD1 and BD2, the pressure levels of the cycle, i.e. the inlet syngas pressure and the final pressure after BD2, and the flow rate of the feed gas.

For the PSA process, the sorbent is a commercial activated carbon, which is particularly suitable for H₂-CO₂ separation. The adsorption related parameters used are those reported in Marx et al. 41.

The feed syngas has the following molar composition:

- CO₂: 38.3%
- H₂: 53.8%
- N₂: 7.4%
- Ar: 0.5%.
In order to determine the energy performance of the adsorption cycle, we consider that the feed syngas is available at 34.1 bar, in line with existing reforming processes, and that only further compression represents an energetic cost for the process. The separated CO\textsubscript{2} stream is designated to storage, and therefore compressed to 120 bar. The energetic cost for CO\textsubscript{2} compression is included, since it is affected by the variation of the blowdown pressure and duration, which determine the CO\textsubscript{2} product. The energy consumption of the CO\textsubscript{2} compressor is calculated with commercial software Aspen Plus \textsuperscript{®}, considering a state of the art compressor configuration.

**Problem formulation**

The objectives considered are the specific energy consumption, and the productivity. The energy consumption, as electricity need, is per mass of produced CO\textsubscript{2}, as the H\textsubscript{2} separation could be carried out with no energy needs when no CO\textsubscript{2} recovery is considered (the syngas is pressurized before the PSA). In terms of separation performance, the PSA optimization is constrained to achieve a CO\textsubscript{2} purity ($\phi_{CO_2}^{\text{PSA}}$) of 85\%, a CO\textsubscript{2} recovery ($r_{CO_2}^{\text{PSA}}$) of 80\% and a H\textsubscript{2} purity ($\phi_{H_2}^{\text{PSA}}$) of 95\%. The resulting optimization problem is formulated as follows in Eq.24:

\[
\begin{align*}
\min & \quad \text{en}^{\text{PSA}}(x) \\
\max & \quad \text{Pr}^{\text{PSA}}(x)
\end{align*}
\]

Subject to:

\[
\begin{align*}
r_{CO_2}^{\text{PSA}} & > r_{CO_2}^{\text{PSA}_{\text{min}}} \\
\phi_{CO_2} & > \phi_{CO_2}^{\text{PSA}_{\text{min}}} \\
\phi_{H_2} & > \phi_{H_2}^{\text{PSA}_{\text{min}}} \\
\text{lb} & \leq x \leq \text{ub}
\end{align*}
\]

Eq.24
\[ x \in \mathbb{R}^n \]

Slightly different parameters are adopted with respect to the TSA case. As reported in Table 8, a lower number of overall function evaluations is allowed, because of the lower number of decision variables (4 vs 8 for the TSA case study), and in order to limit the computational time required to solve the optimization problem, since the convergence time for the PSA model is considerably longer than for the TSA. Moreover, the maximum number of function evaluations for a single Pareto point is also decreased, because a higher number of fixed penalty runs is suggested when dealing with heavily constrained problem as in the case of the PSA case study, which considers three process constraints.

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Adopted value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of workers</td>
<td>8</td>
</tr>
<tr>
<td>Number of worker turns</td>
<td>2</td>
</tr>
<tr>
<td>Number of Pareto points</td>
<td>18</td>
</tr>
<tr>
<td>Maximum number of function evaluation for a fixed penalty parameter ((N_{\text{max}})_k)</td>
<td>100</td>
</tr>
<tr>
<td>Maximum number of function evaluation for a single Pareto point search ((N_{\text{max}})_{\text{tot}})</td>
<td>500</td>
</tr>
</tbody>
</table>

**Results**

The overall computational time required to run the PSA optimization is 24.6 hours. The overall number of function evaluations required is 7,785. On average, each Pareto point search requires 432 function evaluations. The optimization runs leading to points in the high productivity region appear to be more challenging, resulting in a higher number of function evaluations. This is confirmed by the lack of Pareto points retrieved from the look up tables in region close to the maximum productivity (see Figure 11), which might be explained by the additional active constraint in that region of the Pareto curve. The main contribution to the increased complexity of the PSA case study with respect to the TSA is given by the larger number of constraints, and the fast change of the adsorption fronts when pressure is increasing or decreasing. Moreover, the use of look up tables is limited (7% of total function evaluations are retrieved from the look up table, compared to the 12% for the TSA case study): this because the objective functions are more sensible towards the penalty factor (i.e. the shape of the objective functions changes more significantly when increasing the penalty factors), so the algorithm tends to sample different points in successive fixed penalty runs. The sampled points which fail to converge are assigned a suitable suboptimal value but they are not written down into the look up tables.

The Pareto curve obtained for the PSA case study exhibits a clear trade-off between energy cost and productivity, as shown in Figure 11. All the points satisfy purity and recovery constraints. The productivity ranges from 1,025 to 1,712 \(\frac{kg\text{CO}_2}{kg\text{bed} \cdot h \cdot \text{cycle}}\), while the energy cost varies between 0.368 and 0.973 \(\frac{MJ}{kg\text{CO}_2}\). It is worth noting that the energy consumption of the PSA system includes the energy cost of the compressor of the CO\(_2\) produced stream. The selected starting point is close to the Pareto curve; however, it does not meet the constraints of the problem, hence cannot be regarded as optimal solution. The resulting spacing of Pareto points is satisfactory and, interestingly, a discontinuity in the slope of the Pareto curve is identified: in the low productivity region, the productivity can be enhanced with a reduced increase of energy consumption, while after the discontinuity the curve is steeper, meaning that further increasing the productivity becomes more energy expensive. This is because of the inlet stream compressor, whose degree of utilization is decided by the optimization: for low productivity and energy consumption, the syngas feed pressure is kept at 34.1 bar, with no consumptions for the PSA section but only for the CO\(_2\) compression. For high productivity, the optimizer
tends to increase the syngas pressure. In fact, this is equivalent to a superstructure-like approach: if the pressure of the syngas is kept to its lower bound, no syngas compressor is not adopted, while it is when higher syngas pressures are optimal. Accordingly, the discontinuity found for productivity close to $1300 {\frac{kg_{CO_2}}{kg_{bed} h_{cycle}}}$ corresponds to switching on/off the syngas compressor. This is confirmed also in Figure 12 D, where the trend of the feed stream pressure is outlined: low pressure can be found in the low productivity region, high pressure in the high productivity region.

The trends of the decision variables are represented in Figure 12. As mentioned before, the trend for the pressure of the feed stream (Figure 12 D) retraces the behaviour of the Pareto curve. On the other hand, it is difficult to appreciate clear trends for the other decision variables; however MO-MCS identifies a narrow optimal range within the considered decision variables space and adjusts the other decision variables according to the predominant effect of the feed pressure. Such optimal ranges are highlighted with dashed lines in Figure 12 A, B and C. Moreover, we can observe that the constraint on both H$_2$ and CO$_2$ purity kicks when the feed compressor is active, while they are not active in the low productivity region. It is worth noting that the constraint on CO$_2$ recovery is never active, meaning that it is dominated by other constraints (i.e. a solution which meets H$_2$ and CO$_2$ purity constraints, also meets the specification on CO$_2$ recovery). The trend for the syngas flow rate is not reported in Figure 12, since it is constant at equal to the upper bound assumed for the

![Figure 11 Pareto curve resulting from the MO-MCS run. The filled dots are the point returned directly from the algorithm, while the empty circles are the non-dominated points found in the look-up tables after MO-MCS is executed. Note that the selected initial guess (filled triangle) does not meet the constraints of the problem, hence it is not a feasible solution and the associated values of the decision variables are far from the one of the Pareto optimal ones. The objective function values seen by the optimizer (with the first guess penalty factor) is shown with an empty triangle and is considerably distant from the Pareto curve.](image-url)
optimization. Such upper bound is not further extended in order to stay in the validated range of conditions for the adsorption model. The resulting $\text{H}_2$ recovery is minimum for the minimum energy consumption configuration (61.3%), and it increases moving towards higher productivity design up to 80%. The maximum $\text{H}_2$ recovery design corresponds to the maximum productivity one. Overall, the adoption of MO-MCS has allowed to achieve remarkable performance with a rather simple adsorption cycle (in fact, the novel step is the blow down of the first $\text{CO}_2$ product).

Figure 12 Trends of the decision variables for the PSA case study. For each decision variable, the trend is plotted against normalized productivity. A: Time step adsorption; B: Time step blow down 1; C: Time step blow down 2; D: pressure of the feed stream. The empty circles represent optimal solutions were no constraint is active, while filled circles represent points where the constraint on $\text{H}_2$ purity is active, smaller green circles represent points where the constraint on $\text{CO}_2$ purity is active. The constraint on $\text{CO}_2$ productivity is never active.

7 Conclusions

This paper proposes a novel algorithm, named MO-MCS, that was specifically devised to tackle the optimization of adsorption processes with a black-box approach. The algorithm is based on MCS which is a model-based derivative-free algorithm with known convergence results. MCS has been extended to deal with constrained bi-objective optimization problems with the normalized normal constraints method, which enables the calculation of Pareto sets that are (i) well-spaced, and (ii) non-smooth and non-convex. In order to save computationally expensive function evaluations, look up tables has been used to retrieve the objective values of already sampled decision variable vectors, and the calculation of Pareto points exploits parallel computing on multiple workers.

MO-MCS was benchmarked against three other multiobjective optimization algorithms for a TSA test case. MO-MCS performs better than all such algorithms, returning a Pareto curve which is 6% more extended than the second best algorithm, closer to the best Pareto envelope (84% improvement with respect to MOPSO) and well-spaced Pareto points. The computational time required for a single run of MO-MCS is about 35% longer than the one required for NSGA-II and MOPSO because during the search of the extreme Pareto points MO-MCS does not fully exploits parallel computing. On the other hand, the stochastic nature of NSGAII and MOPSO calls for multiple runs of the problem, which makes MO-MCS more efficient also from a computational point of view.
The application of MO-MCS to relevant TSA and PSA case studies shows that the proposed optimization algorithm can be used as an effective and flexible tool to guide the design of adsorption based processes. In fact, MO-MCS can assess the overall potential of an adsorption process in terms of productivity and energy potential, and moreover it is able to return information on the influence of each considered decision variable on the process, thus identifying possible trends as well as its optimal range.

8 Acknowledgments

Federico Capra acknowledges RSE (Ricerca Sistema Energetico) and LEAP (Laboratorio Energia & Ambiente Piacenza) for supporting his research activity.

Supporting Information. The main equations describing the adsorption model adopted in this work are available in the supporting information material. This information is available free of charge via the Internet at http://pubs.acs.org/

Nomenclature

\( b \) Sips constant [1/bar] as a function of temperature
\( c \) fluid phase concentration, [mol/m\(^3\)]
\( C_{ads} \) heat capacity of the adsorbed phase, [J/(K kg)]
\( C_g \) heat capacity of the gas, [J/(K m\(^3\))] 
\( C_s \) heat capacity of the solid, [J/(K kg)]
\( d_p \) particle diameter for packing material [m]
\( e \) energy performance index
\( e_{n} \) specific energy cost [MJ/kg]
\( e_{x} \) specific exergy cost [MJ/kg]
\( f \) objective function
\( g \) constraint function
\( h \) normalized normal constraint
\( h_L \) heat transfer coefficient (lumping column inside +wall), [J/(m\(^2\) s K)]
\( l \) pareto curve extension
\( lb \) lower bound vector
\( N \) utopia vector
\( N \) normal vector
\( N_{max \ k} \) maximum number of function evaluation for a fixed penalty run
\( N_{max \ tot} \) maximum number of function evaluation for a single Pareto point search
\( P \) pressure
\( Pr \) productivity
\( q \) solid phase concentration [mol/kg]
\( q_{\infty \ ref} \) Sips isotherm parameter [mol/kg]
\( r \) recovery
\( R \) column radius (reactor model)
\( \mathcal{R} \) universal gas constant
\( sp \) exponent in Sips isotherm [–]
\( SW \) sweep
\( T \) temperature
\( t \) time coordinate (reactor model)
\( T_{amb} \) ambient temperature [K]
\( u \) superficial gas velocity, [m/s]
\( ub \) lower bound vector
\( x \) decision variables vector
\( X \) utopia line points
\( y \) mole fraction [-]
\( z \) spatial coordinate (reactor model)
\( \delta \) normalized increment
\( \Delta H \) isosteric heat of adsorption [J/mol]
\( \varepsilon_b \) bed void fraction
\( \varepsilon_t \) overall void fraction, [-]
\( \mu \) penalty factor
\( \rho_b \) fluid phase density, [kg/m\(^3\)]
\( \rho_p \) particle density, [kg/m\(^3\)]
\( \sigma \) spacing constant
\( \tau \) penalty update factor
\( \phi \) purity

**Acronyms**

- BDF backward differentiation formula
- CSS cyclic Steady State
- DAE differential Algebraic Equations
- DMS direct Multi Search
- LDF linear Driving Force
- MCS multilevel Coordinate Search
- MO-MCS multi Objective Multilevel Coordinate Search
- MOPSO multi objective Particle Swarm
- NBI normal-boundary intersection
- NG natural gas
- NNC normalized normal constraint
- NSGA II non sorting genetic algorithm II
- PDE partial derivative equations
- PSA pressure swing adsorption
- TSA temperature swing adsorption
- VPSA vacuum pressure swing
- WS weighted sum

**Subscripts/superscripts**

- \( A_1 \) Pareto extreme Point 1
- \( A_2 \) Pareto extreme Point 2
- \( b \) Bulk phase
- \( C \) Constraint
- \( PP \) Pareto Points
- \( s \) Solid phase
- \( W \) Worker
- \( WT \) Worker Turns
References


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